



# CISPA

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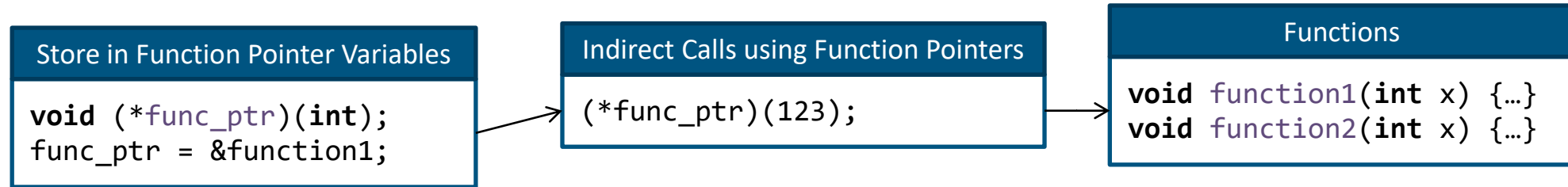
# TyPRO: Forward CFI for C-Style Indirect Function Calls Using Type Propagation

**Markus Bauer, Ilya Grishchenko, Christian Rossow**

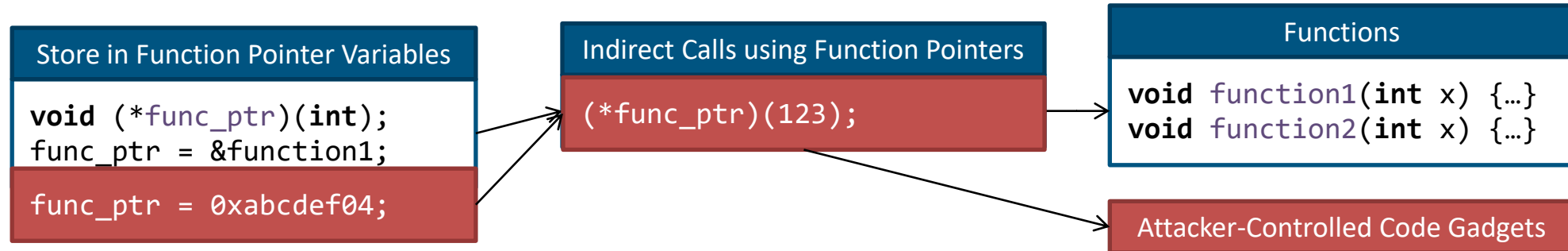
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Code: <https://github.com/typro-type-propagation/TyPro-CFI>

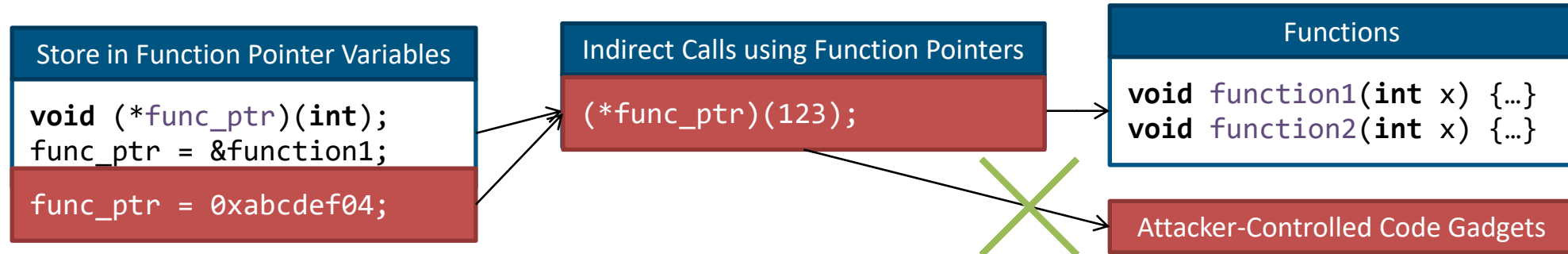
# Function Pointers and Indirect Calls in C



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# Control Flow Integrity (CFI) for Indirect Calls



- Indirect Calls should only call legitimate functions
- Control Flow Integrity Schemes:
  - Build **Target Sets** for each indirect call
  - Ensure only valid targets can be called

TargetSet = {function1, function2}

# Control Flow Integrity (CFI) in Practice


## functional

Control Flow Guard  
(Microsoft)

- Each function is a target for each indirect call
- **Risky overapproximation**  
(possibly >1000 targets)

## most possibly secure

Clang CFI

- A function is a target if function type == pointer type
- **Breaks applications** 

## Problems of Strict CFI

### Unmatched Argument Types (lighttpd, hmmmer)

```
long (*filter_ptr)(void *, data *);
long filter_func (plugin *p, data *d) {}

int (*compare_ptr)(const void *, const void *);
int compare_func (long a, long b) {}
```

### Variadic Function Types (gcc, nginx)

```
rtx (*gen_ptr)(rtx, ...);
rtx gen_func (rtx a, rtx b, rtx c) {}

handler_t (*mod_ptr) (...);
handler_t mod_deflate(request *r, void *data) {}
```

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## functional and secure

TyPro

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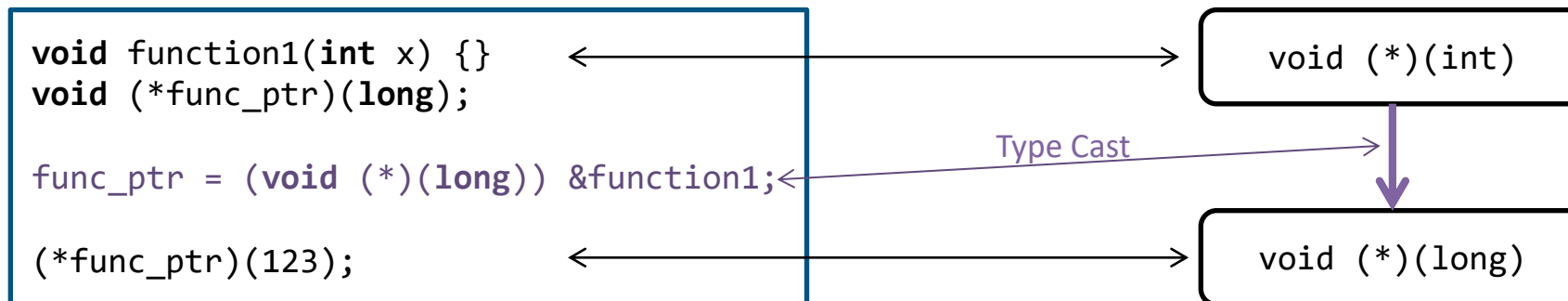
# TYPRO: Type Propagation

- C compiler prototype based on Clang / LLVM 10
- Protects indirect calls from memory corruption vulnerabilities
- Protects existing software *without modifications*
- Features:
  - Multiple architectures (x86, ARM, MIPS)
  - 100% protected binaries (using musl libc)
  - Dynamic loading of additional code
  - No performance impact



# Our Approach #1: Track C Type Casts

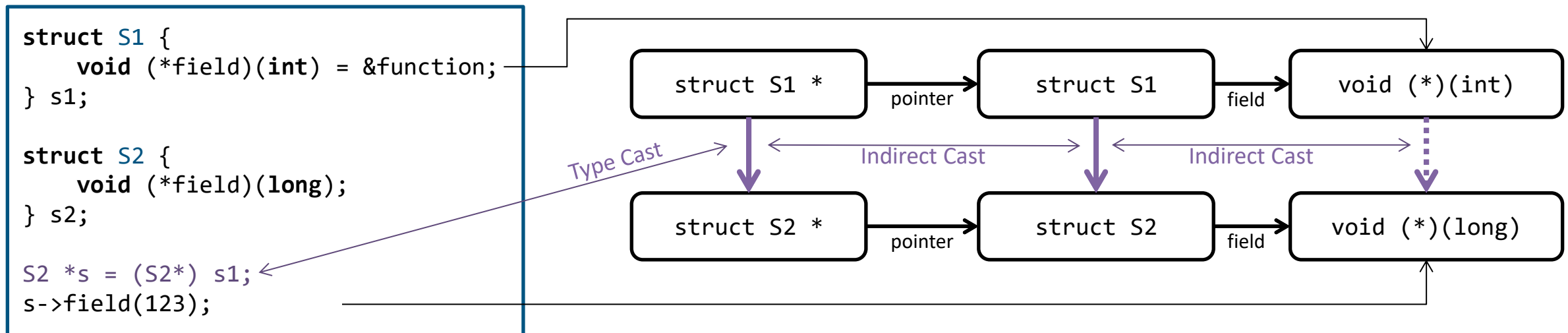
- Types of function pointers change is explicit: Type Casts
- Collect all type casts from C source
- Track reachability between function and indirect call types





# Our Approach #2: Indirect Type Casts

- Function pointer types can be part of composed types (structs, pointers, ...)
- We must track implicit changes (from composed types)



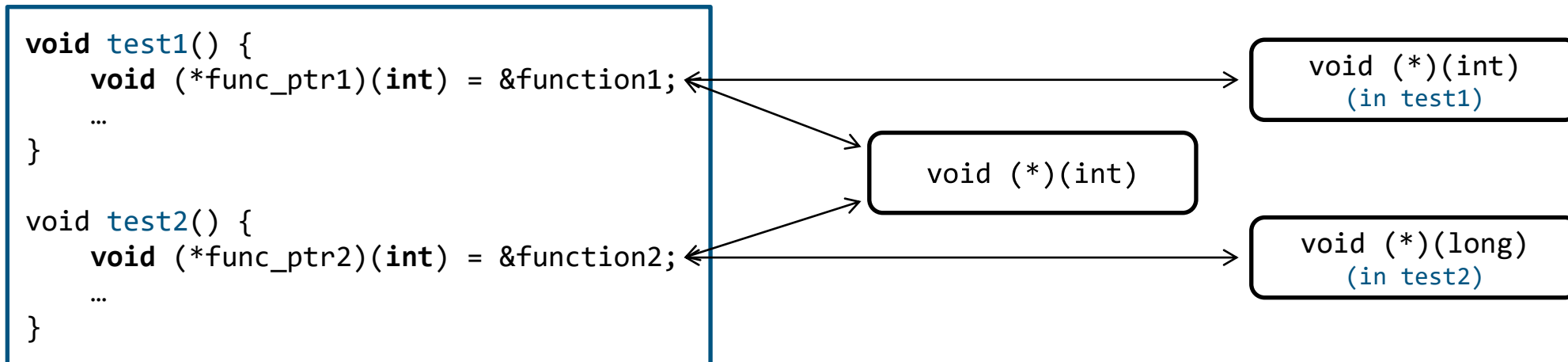
# Our Approach #3: Track Types per Function

- One general cast in one function could affect all target sets

For example: casts from/to void\*

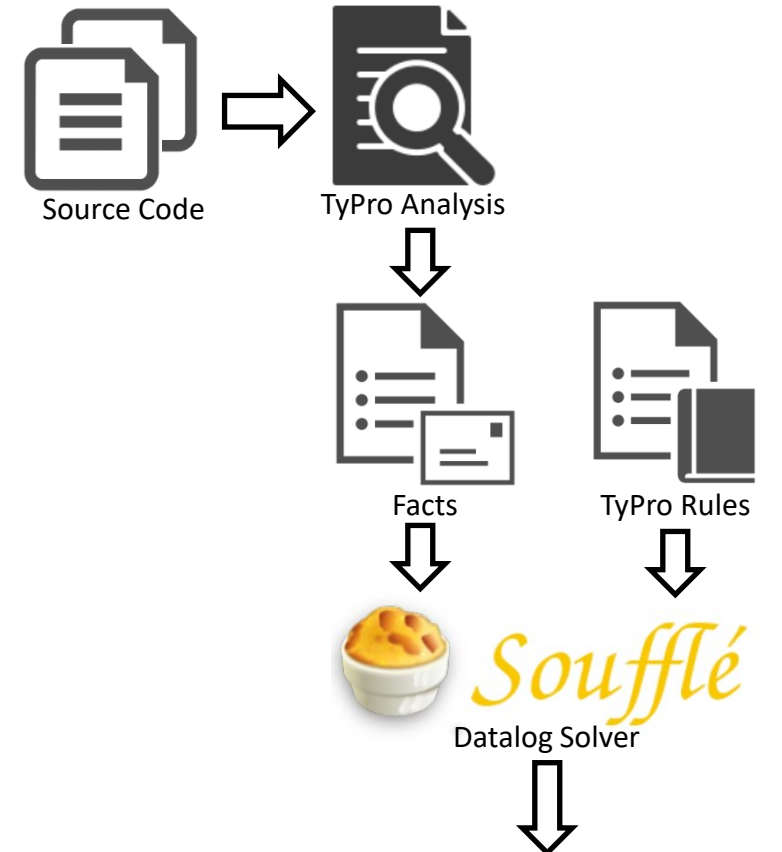
→ We consider types and casts *per function*

→ (and the type propagation between functions)



# TyPRO: Type Propagation

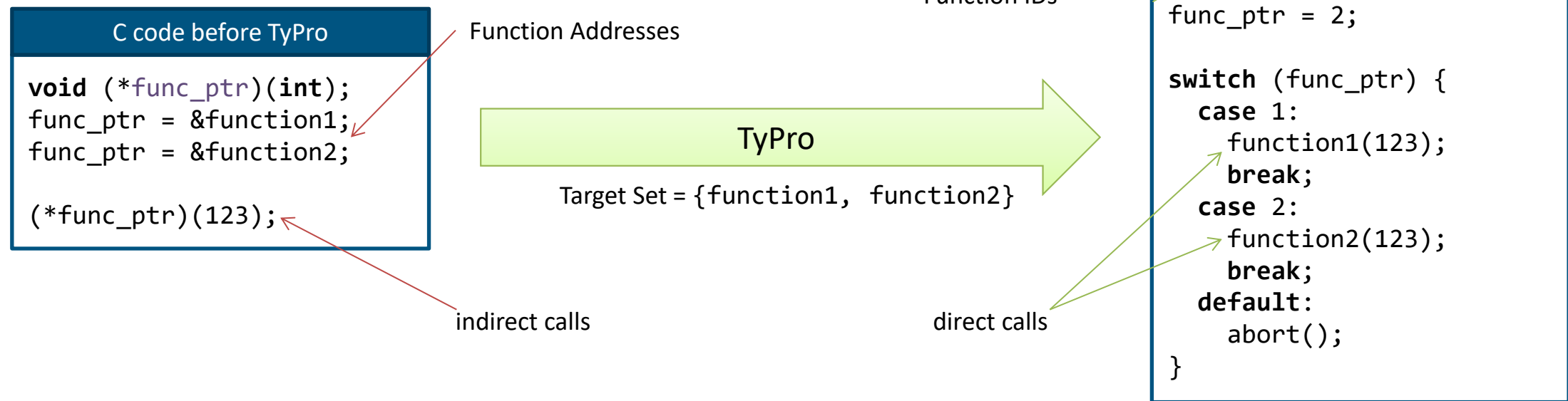
- Uses a formal, rule-based system to follow type casts
- Collect types and casts during compilation (“facts”)
- Compute indirect casts and reachability during linking  
(using pre-defined “rules”)
- Compute target sets for all indirect calls



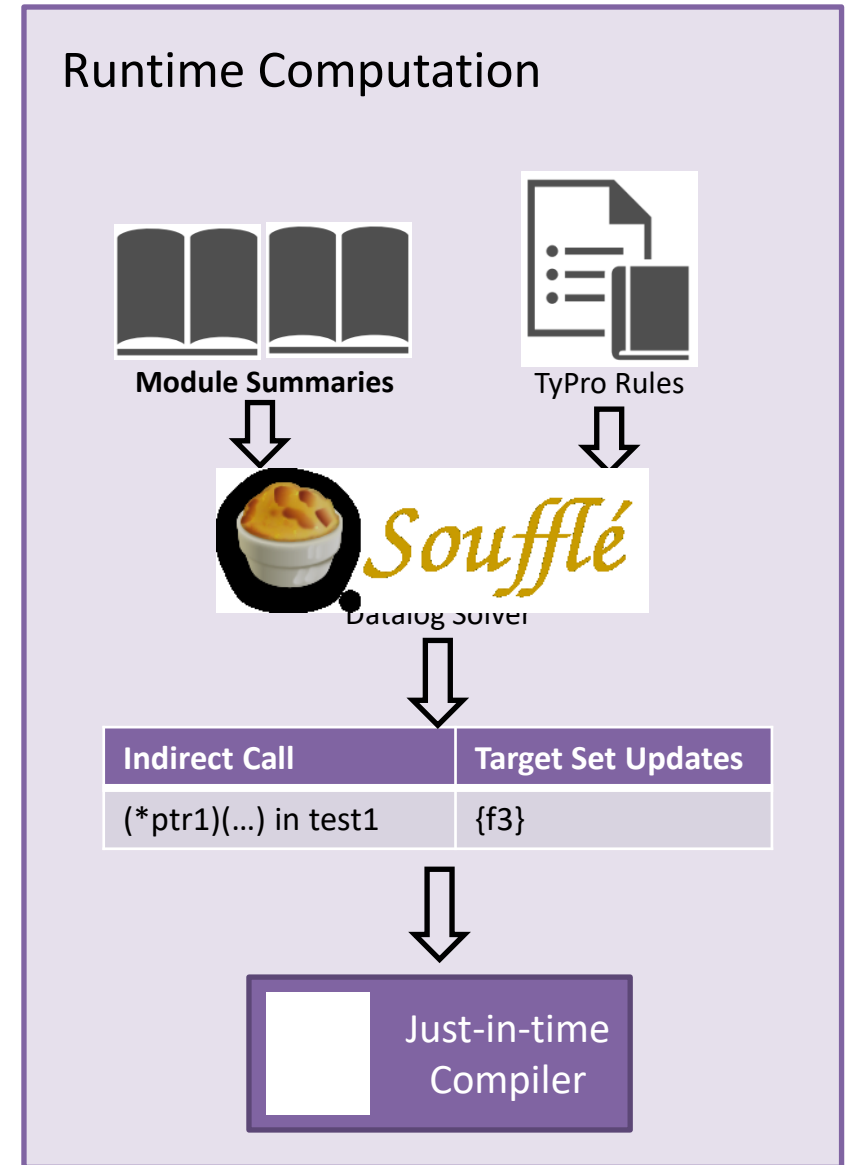
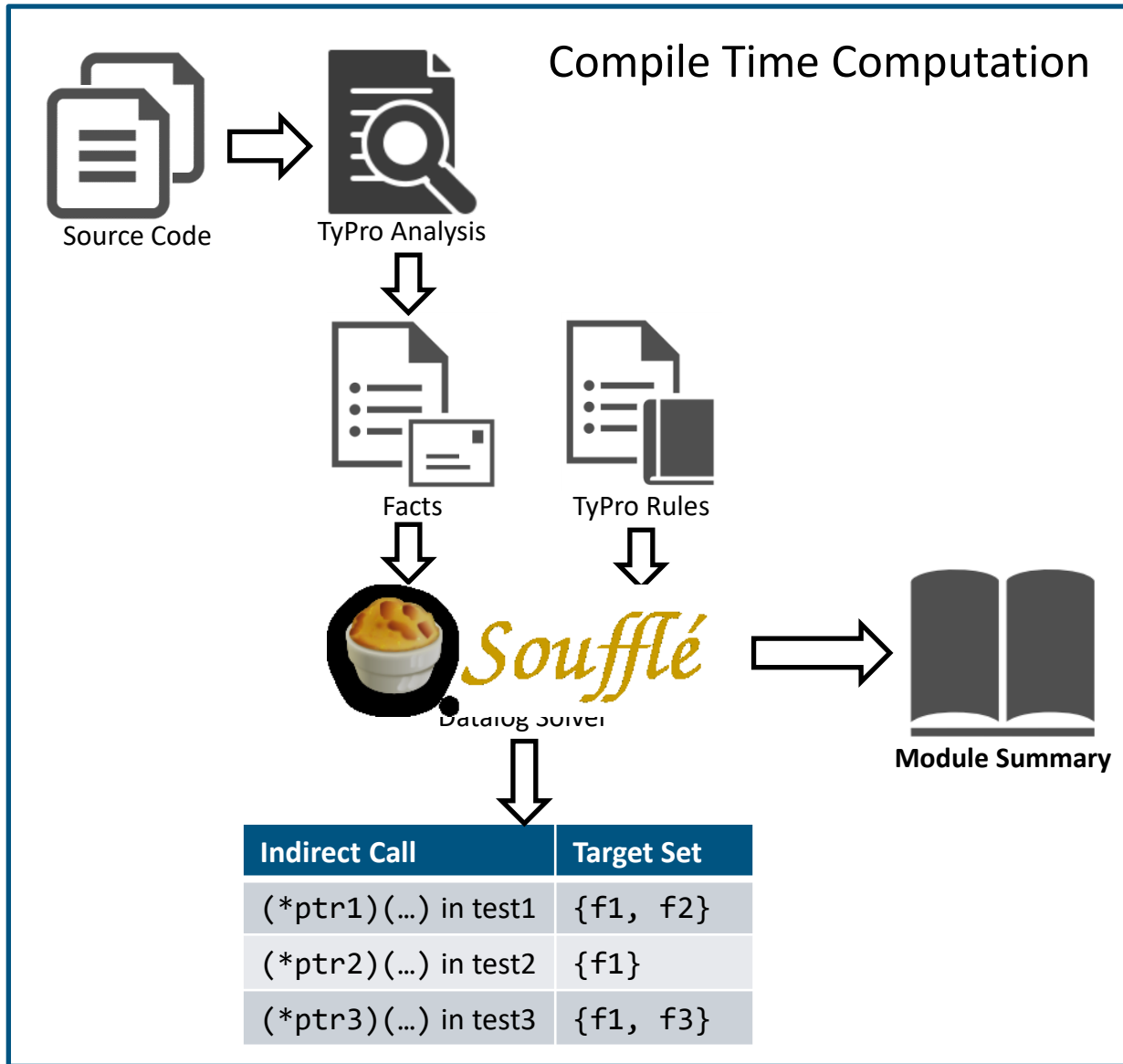
Indirect Call	Target Set
(*ptr1)(...) in test1	{f1, f2}
(*ptr2)(...) in test2	{f1}
(*ptr3)(...) in test3	{f1, f3}

# Enforcing Target Sets

- Replace function pointers by numbers (function ID)
  - Replace indirect calls with `switch` over safe, direct calls
- no unsafe indirect call remains



# Dynamic Linking – Adding Code at Runtime



# Evaluation of Type Propagation

**functional**

**Control Flow Guard**  
(Microsoft)

**functional and secure**

**TyPro**

**most possibly secure**

**Clang CFI**

**SPEC CPU 2006 is:**  
perl, bzip2, gcc, milc,  
gobmk, hmmer, sjeng,  
h264ref, sphinx3

**SPEC CPU**  
up to 1786 targets

**SPEC CPU**  
removed ~79.0%  
of all targets

**SPEC CPU**  
removed ~79.2%  
of all targets  
**2/9 applications break**

**Real-World Applications:**



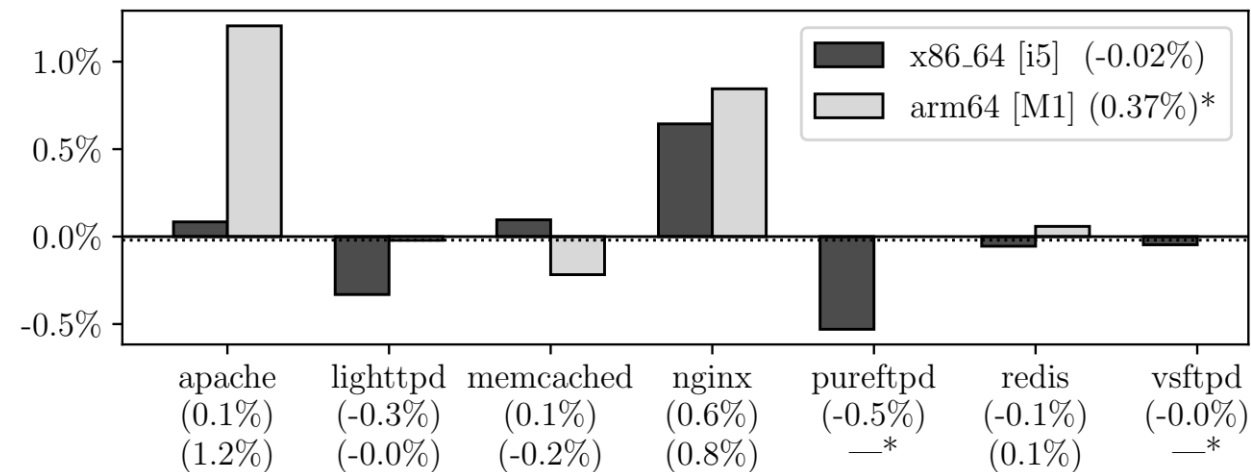
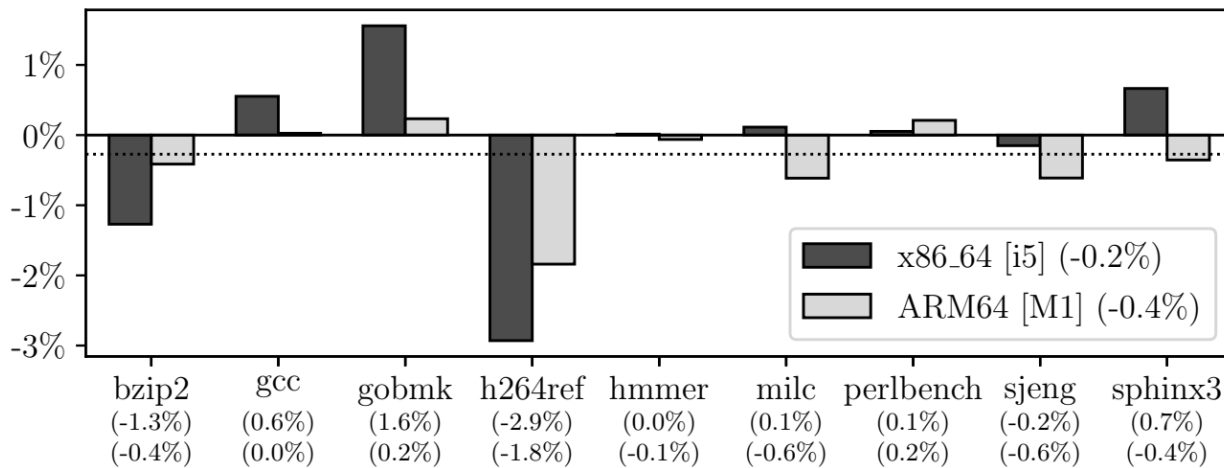
**Real-World Applications**  
up to 2267 targets

**Real-World Applications**  
removed ~95.2%  
of all targets

**Real-World Applications**  
removed ~97.6%  
of all targets  
**4/7 applications break**

# TYPRO: Performance

- On average, programs get (negligibly) faster
  - No checks necessary
  - More suitable compiler optimizations
- On average, programs get ~9% larger



# Summary / Questions?

- Type-based analysis following casts
- Compute small yet correct target sets
- Rewrite indirect into direct calls
- Compatible, secure, and fast

```
C code after TyPro

void (*func_ptr)(int);
func_ptr = 1; // Function ID

switch (func_ptr) {
  case 1: function1(123); break;
  case 2: function2(123); break;
  default: abort();
}
```

