

THE CHATTY SENSOR: A PROVABLY-COVERT CHANNEL IN CYBER PHYSICAL SYSTEMS

Amir Herzberg

amir.herzberg@uconn.edu

University of Connecticut

Storrs, USA

[Yehonatan Kfir](#)

yehonatank@gmail.com

Bar-Ilan University

Ramat-Gan, Israel

INTRODUCTION

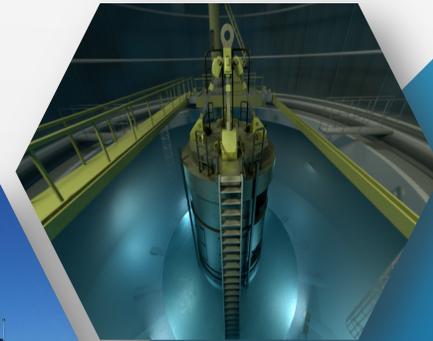
- ▶ **Cyber Physical Systems (CPS)** - Smart systems that include networks of physical and computational components, all aimed to governed a physical process.
- ▶ **Examples:** Nuclear Plants, Power Generations, Water Plant, Transportations.

- ▶ Critical for our life
- ▶ Built from large number of devices:

Sensors, Actuators, Controllers...

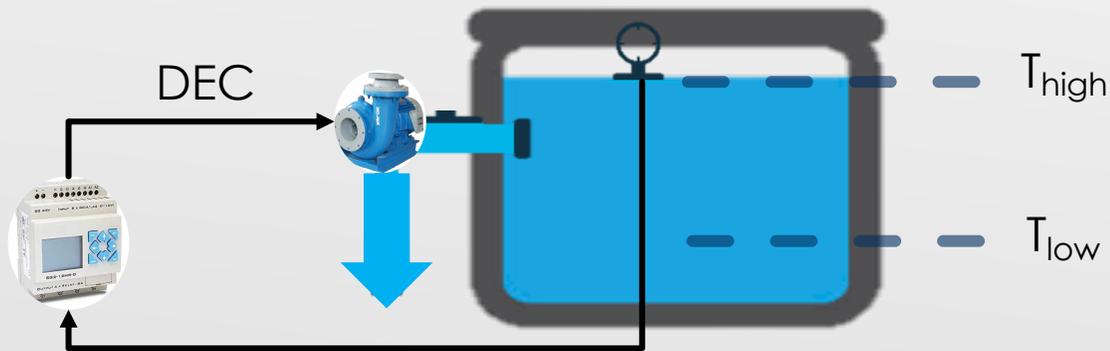


- ▶ Operating in *Feedback Control Loops*



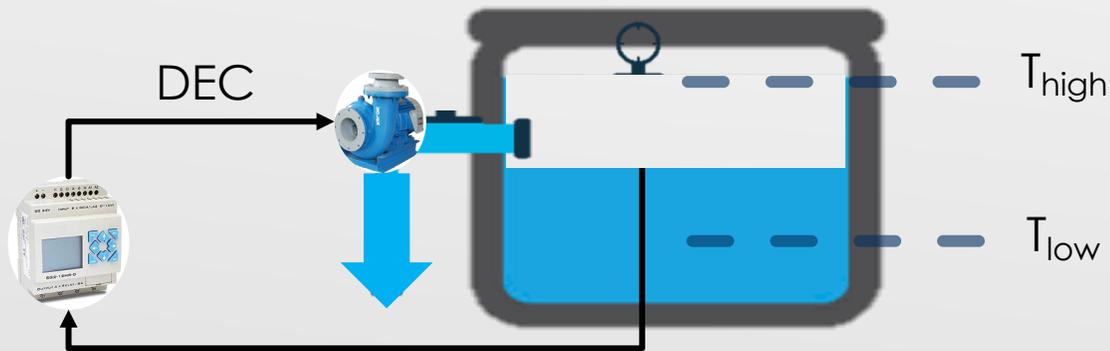
INTRODUCTION: FEEDBACK CONTROL LOOP

- ▶ Feedback control loops are the main method used to stabilize physical values in CPS.
- ▶ Threshold-controller
 - ▶ Actuator with two possible commands to increase / decrease the physical value: *INC* / *DEC*
 - ▶ Two thresholds: T_{high} , T_{low}
- ▶ When the sensor measurements reach T_{high} / T_{low} , the controller changes its output to decrease / increase the signal.



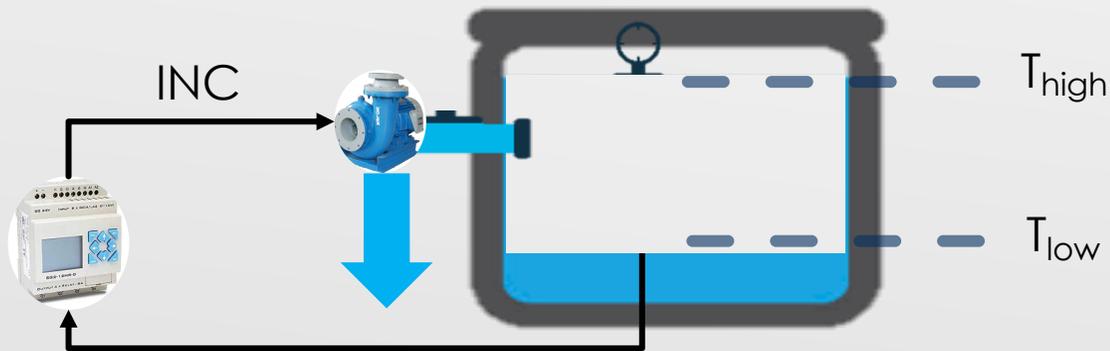
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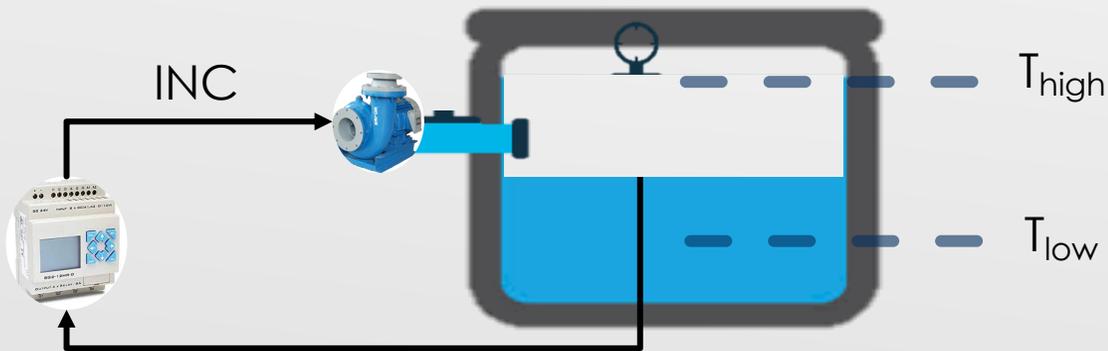
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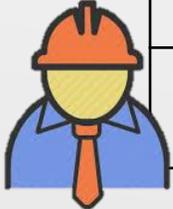
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- ▶ Widely used in: phase controller, current limiter, pH controllers.

INTRODUCTION: DEVICE SELECTION

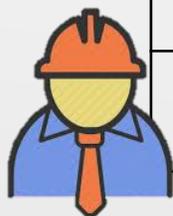
- ▶ Devices are chosen based on **sufficient specification** and **lowest cost**.



	Device A 	Device B 
Specification	 High Quality <small>SPEC</small>	 Sufficient Quality <small>SPEC</small>
Price	 Expensive	 Cheap

INTRODUCTION: DEVICE SELECTION

► Devices are chosen based on **sufficient specification** and **lowest cost**.



	Device A 	Device B 	Malicious 
Specification	 High Quality <small>SPEC</small>	 Sufficient Quality <small>SPEC</small>	 Sufficient Quality <small>SPEC</small>
Price	 Expensive	 Cheap	 Cheaper

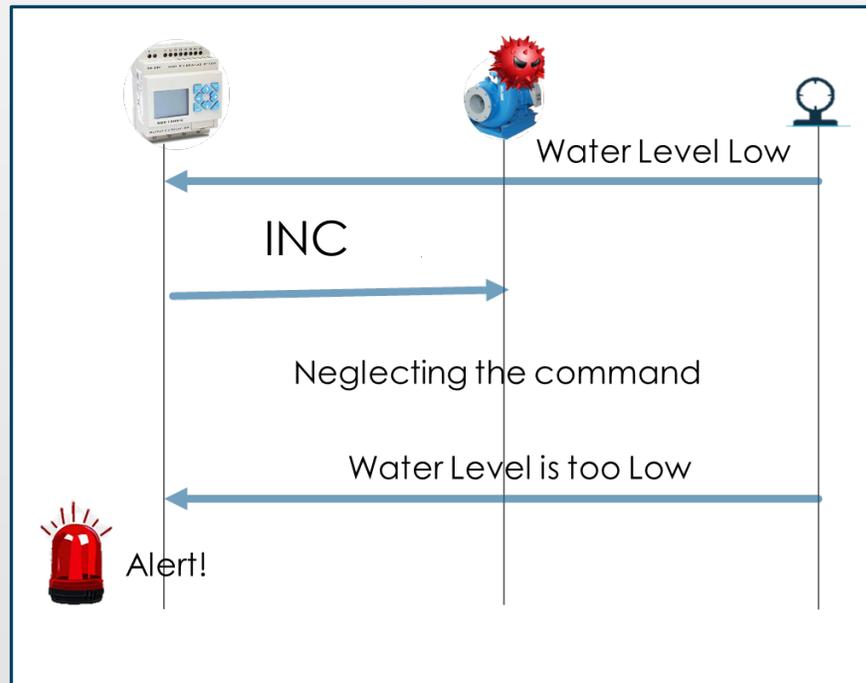
► Supply Chain Attack:

- Attacker offers a cheaper device, with sufficient specification.
- OR: Attacker replaces benign devices, with malicious one.

► **Attacker Goal:** To cause damage, by deploying its own malicious device.

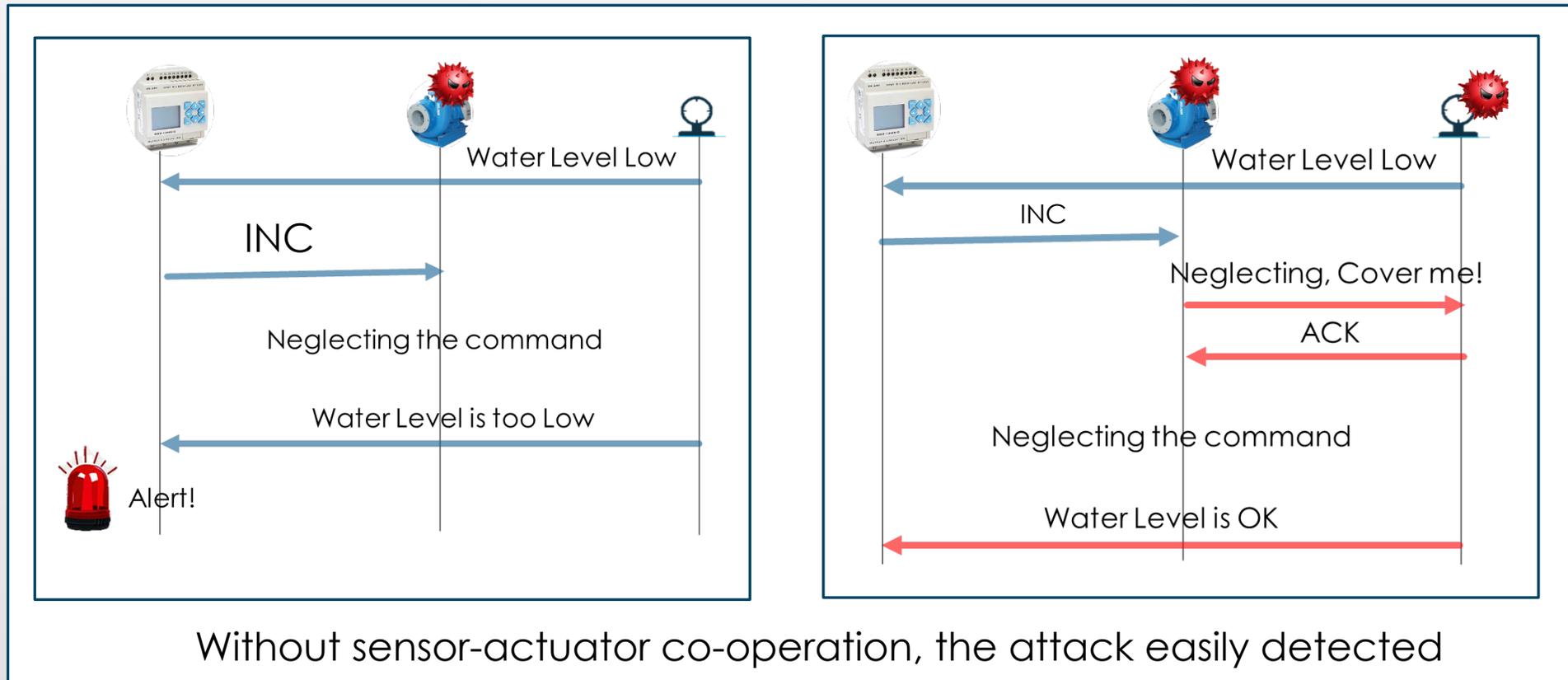
ATTACKER CHALLENGE - 1

- ▶ Successful, stealthy attack requires communication
 - ▶ e.g. from corrupt sensor to corrupt actuator:



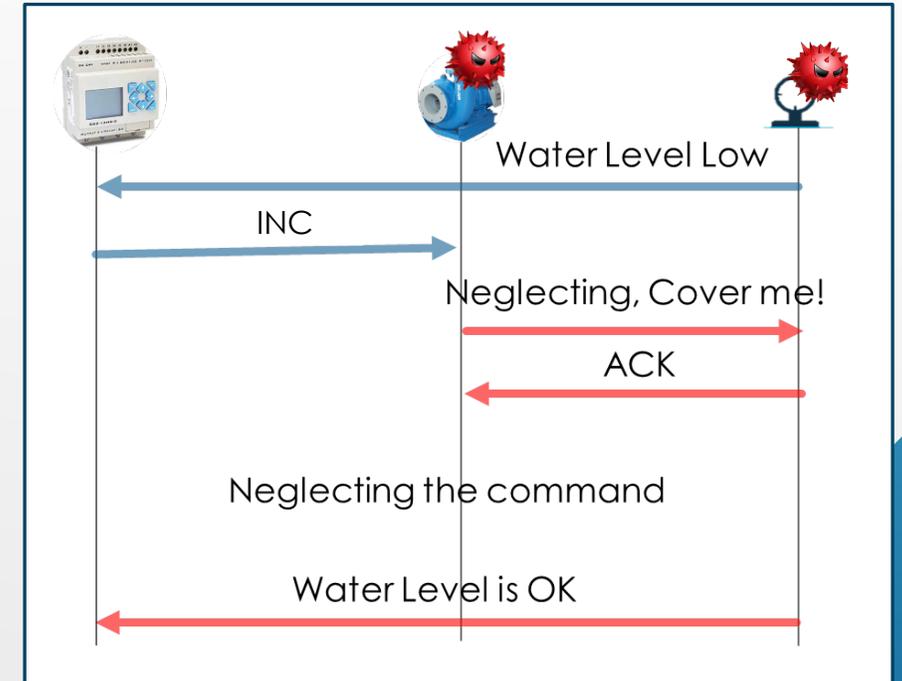
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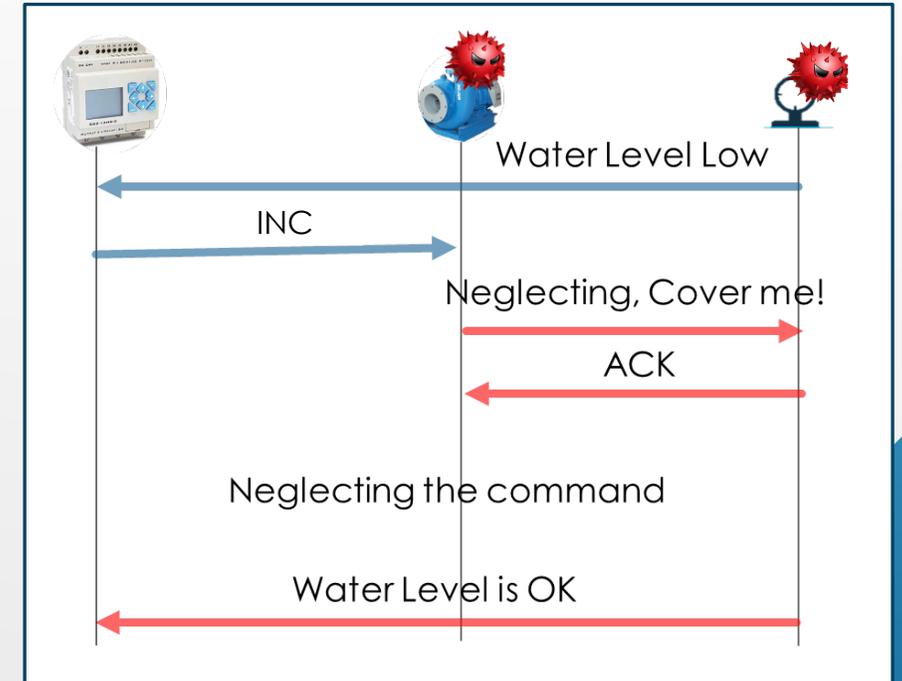
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- ▶ How to communicate **between** malicious devices?
 - ▶ Sensor to Actuator (S2A) – This work.
 - ▶ Actuator to Sensor (A2S) – Prev. work.



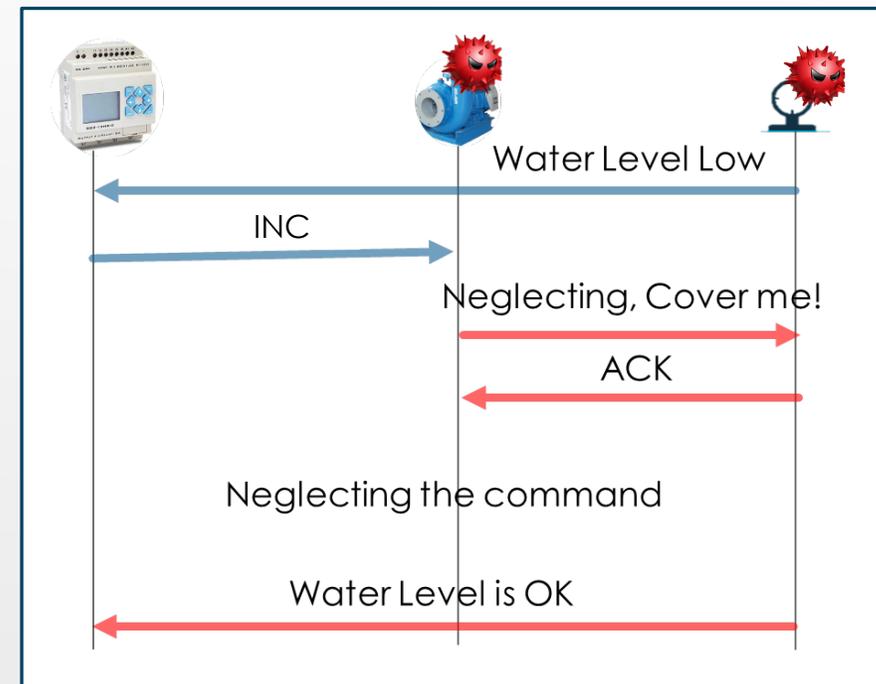
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- ▶ How to avoid detection?



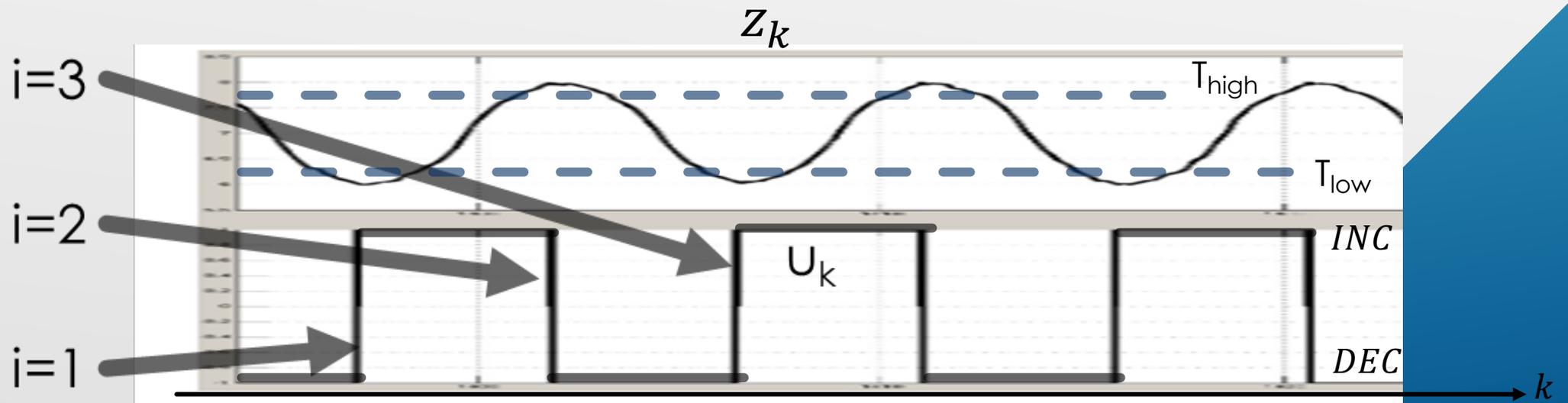
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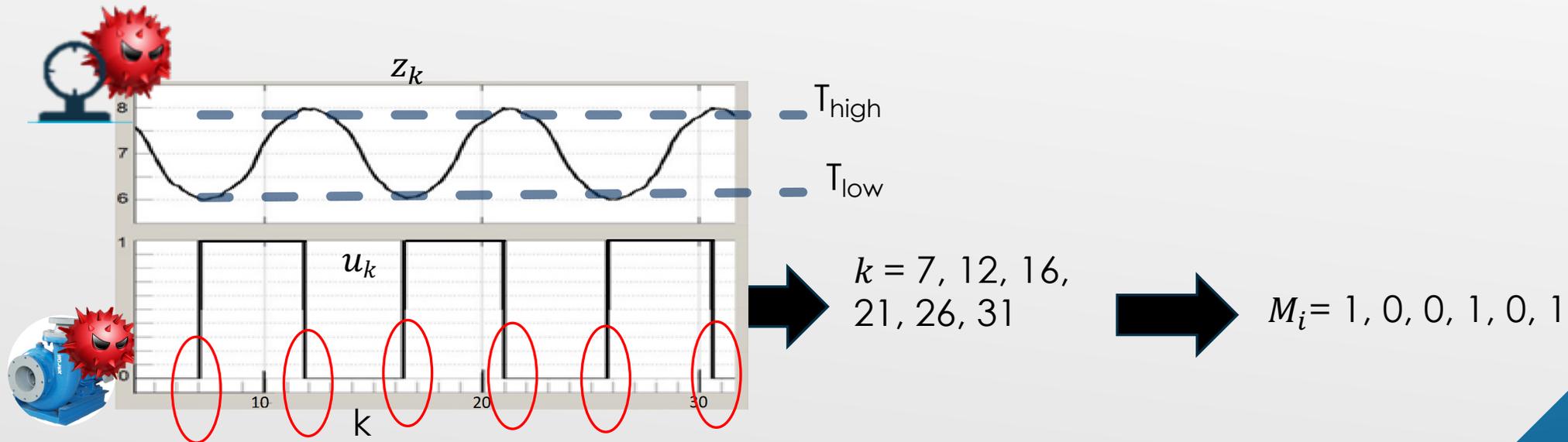
CHATTY-SENSOR COMMUNICATION METHOD

- ▶ For any time-step k , the **sensor** reports z_k .
- ▶ The process value continuously iterates and pass the thresholds: $T_{\text{high}}, T_{\text{low}}$
- ▶ Whenever z_k passes a threshold, the controller **switches** the command $u_k \in \{INC, DEC\}$ to the **actuator**.
- ▶ We denote the i^{th} transition of the actuator's output by i .



CHATTY-SENSOR COMMUNICATION METHOD

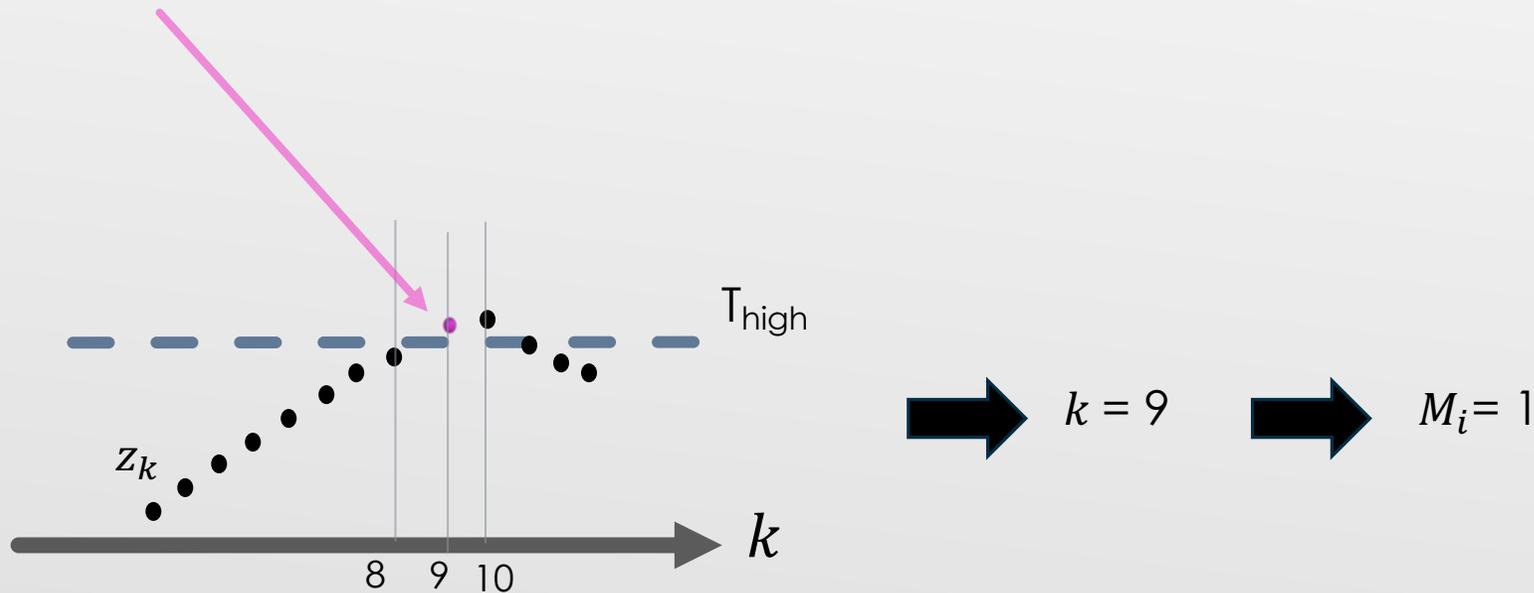
- ▶ Sensor encodes covert bits of information, on the **parity of the transition time-steps**:
 - ▶ Transition at even / odd times will signal bit 0/1.



Assumption: Sensor and Actuator have a parity-synchronized clocks.

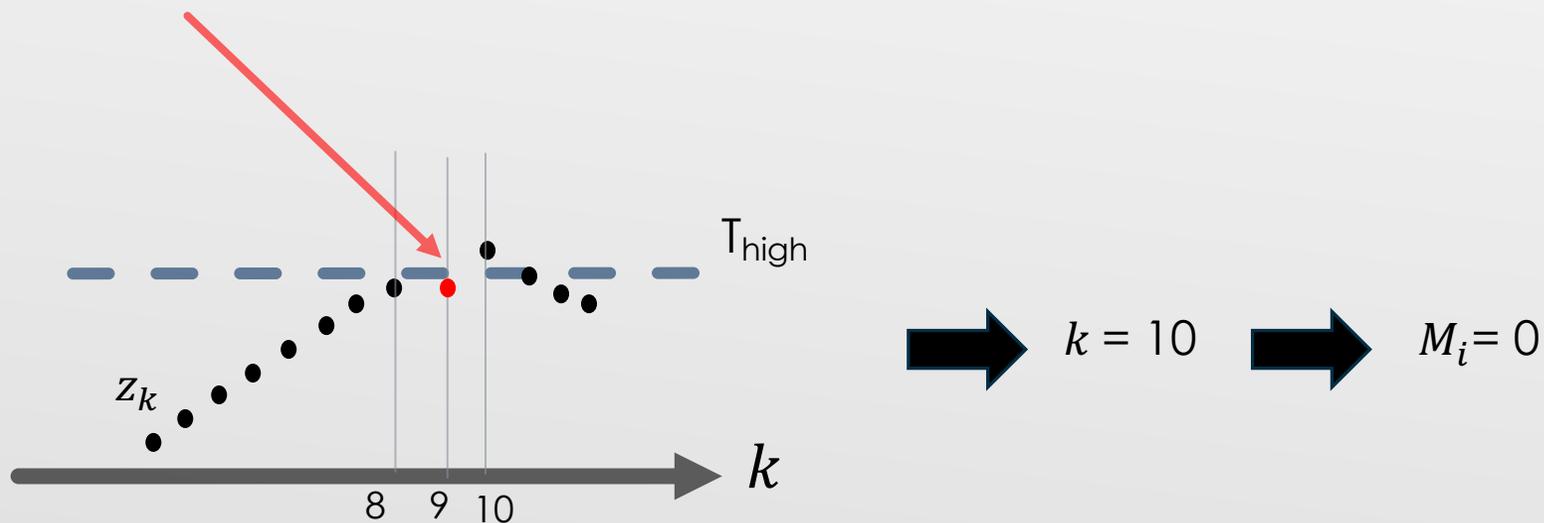
CHATTY-SENSOR COMMUNICATION METHOD

- ▶ Chatty-sensor influences the transition time-step.
 - ▶ Decreasing / increasing the reported value.
- ▶ For example:
 - ▶ Transition about to happen at $k=9$ -> but should be at $k=10$.



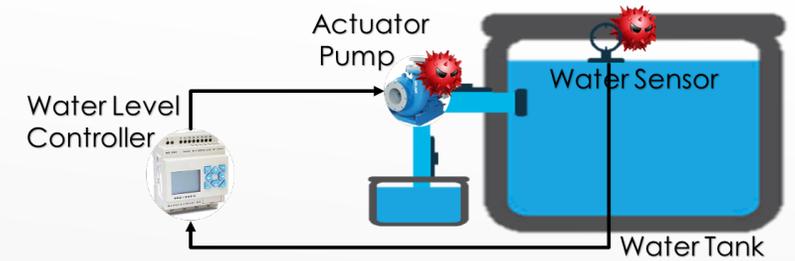
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ATTACKER CHALLENGES

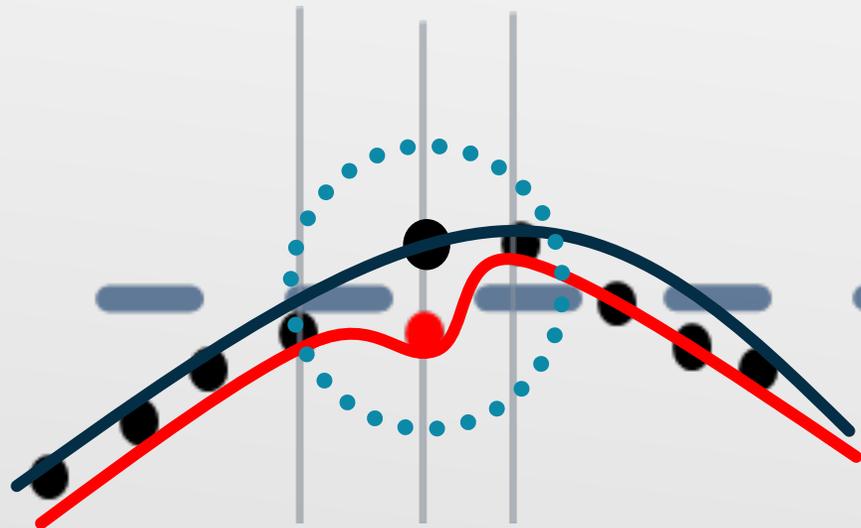
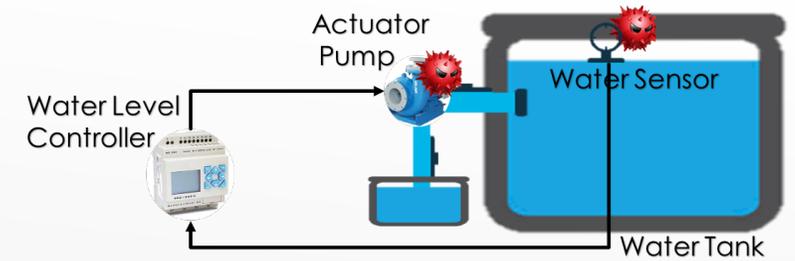
► How to communicate **between** malicious devices?



Transition Parity

ATTACKER CHALLENGES

► How to communicate **between** malicious devices?



— Chatty-sensor

— Benign sensor

Transition Parity

Creates Anomaly in the CPS behavior...

ATTACKER CHALLENGE - 2

- ▶ A lot of works on anomalies detections in CPS.

- ▶ Communication Network Anomalies:

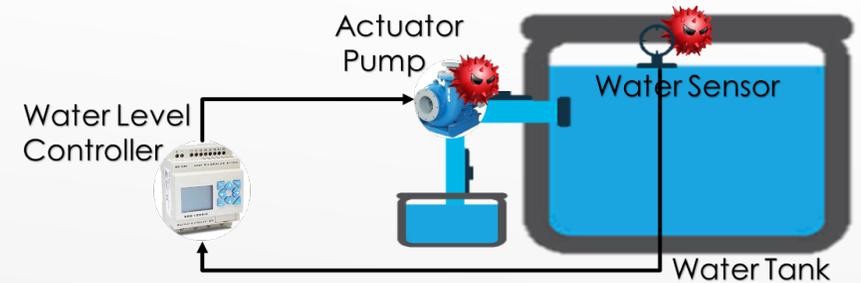
- ▶ For example (one of many):

- ▶ Kleinmann, Amit, and Avishai Wool. "Accurate modeling of the siemens s7 scada protocol for intrusion detection and digital forensics.", 2014.

- ▶ Physical Anomalies – malicious sensor reporting / malfunctioning actuator

- ▶ For example (one of many):

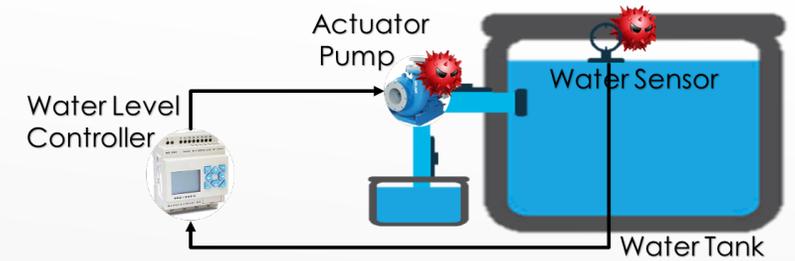
- ▶ Urbina, David I., et al. "Limiting the impact of stealthy attacks on industrial control systems.", 2016.



ATTACKER CHALLENGES

► How to communicate **between** malicious devices?

► How to avoid detection?



Transition Parity

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ATTACKER CHALLENGES

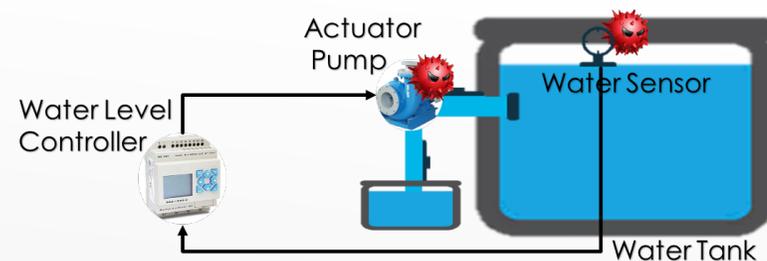
► How to communicate **between** malicious devices?

► **How to avoid detection?**

Covert Channel

Transition Parity

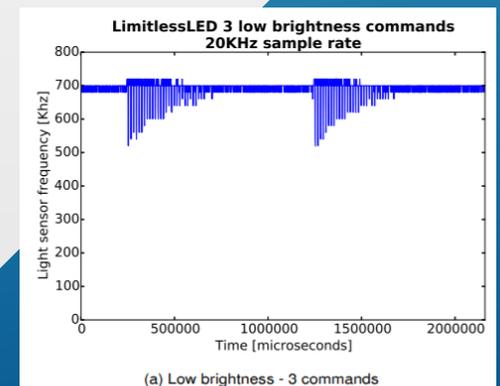
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COVERT CHANNELS

- ▶ “**Covert**” - using some “**unmonitored**” channels
 - ▶ Encoding information using light brightness (“Extended functionality attacks on IoT devices: The case of smart lights“, Shamir et. al. 2016)
 - ▶ Packet headers (“Embedding Covert Channels into TCP/IP”, Murdoch et. al. , 2005)
 - ▶ Acoustic emissions of a motor (“Process-aware covert channels using physical instrumentation in cyber-physical systems”, Krishnamurthy et. al. 2018)
 - ▶ And more...
- ▶ Monitoring the “**unmonitored**” property, reveals the communication channel.

Eyal Ronen and Adi Shamir. Extended functionality attacks on IoT devices: The case of smart lights. In 2016 IEEE European Symposium on Security and Privacy (EuroS&P), pages 3–12. IEEE, 2016



PROVABLE COVERT CHANNELS

▶ “Provable-Covert” –

- ▶ No secret property
- ▶ Proving that it is impossible to detect the channel (under well defined assumptions)

$$\Pr(D(\text{🔍🦠}) = \text{Mal.}) \approx \Pr(D(\text{🔍🕒}) = \text{Mal.})$$

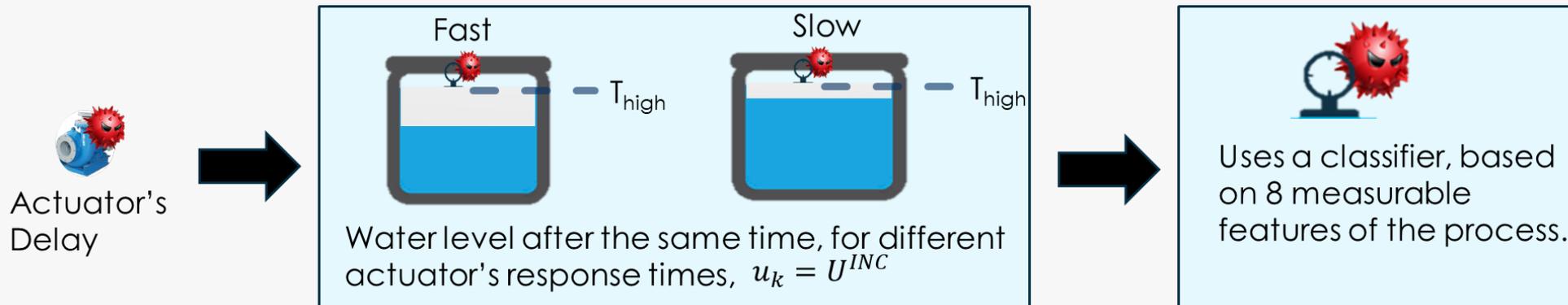
PROVABLE COVERT CHANNELS

► **IT Networks:** Provable channels were presented in the past:

► Liu, Yali, et al. "Robust and undetectable steganographic timing channels for iid traffic.", 2010.

► **CPS Provable Covert Channel:**

Herzberg, Amir, and Yehonatan Kfir. "The Leaky Actuator: A Provably-covert Channel in Cyber Physical Systems." *Proceedings of the ACM Workshop on Cyber-Physical Systems Security & Privacy*. ACM, 2019.



CHATTY-SENSOR COVERT CHANNEL



► How to (provably) avoid detection?

► The provably-covert channel is based on **two basic observations** about sensors:

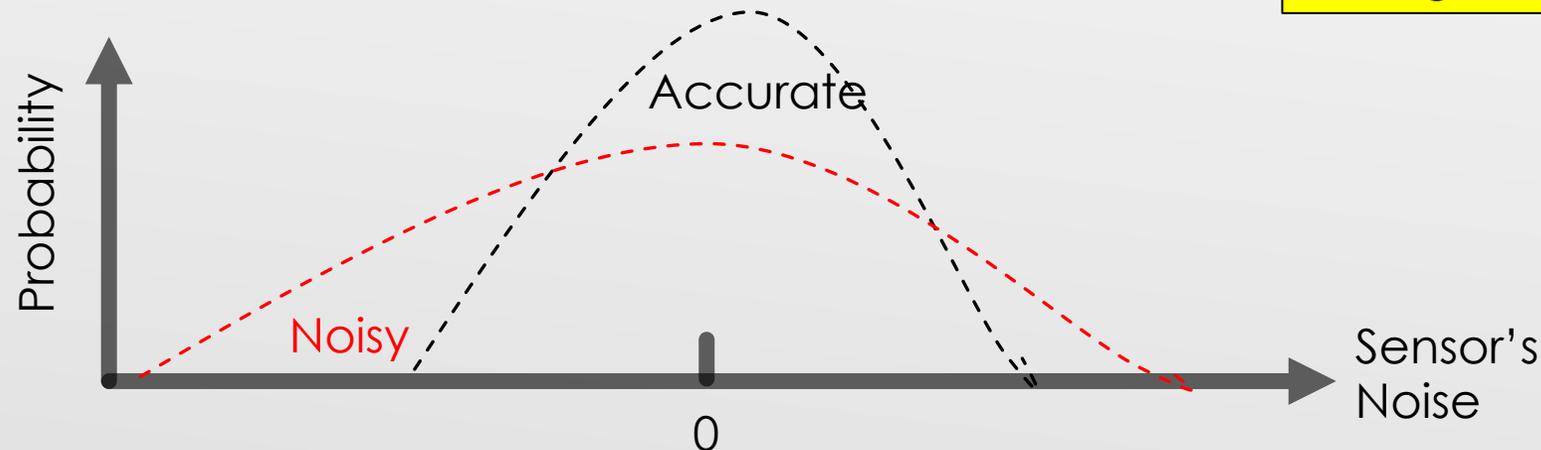
► The **reported measurements has a random** noise, derived from some (known) distribution.

► There are different **benign** types of sensors in the market:

► Accurate (narrow noise distribution)

► Noisy (wide noise distribution).

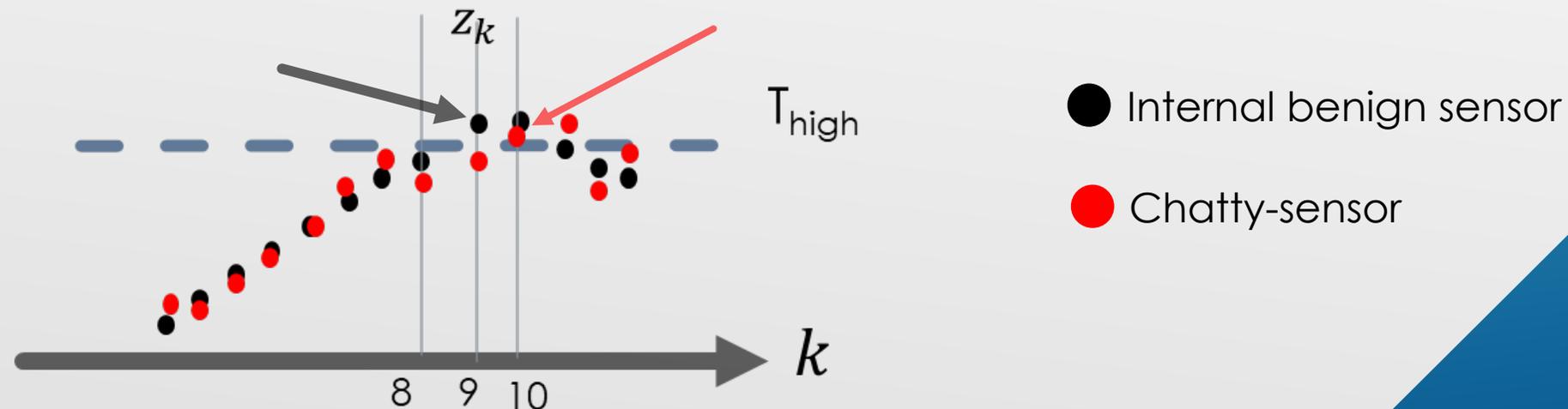
- Adding Noise at all time-steps
- Make sure to add positive / negative noise at the transition time



CHATTY-SENSOR COVERT CHANNEL



- ▶ The Chatty-sensor uses an **internal accurate sensor** to measure the process.
- ▶ At time-steps with transition: the chatty-sensor **chooses** whether to add **positive or negative noise** to the internal sensor.
- ▶ All the other time-steps: The Chatty-sensor **randomly chooses positive or negative noise** to add.



CHATTY-SENSOR COVERT CHANNEL



► **Problem [Encoding]:** Channel is noisy...

► **Solution:** Error Correction Code. Sending encoded message \mathbf{m} , $m = \text{ECC}(M)$.

$$\Pr(D(\text{🕒}) = \text{Mal.}) \approx \Pr(D(\text{👤}) = \text{Mal.})$$

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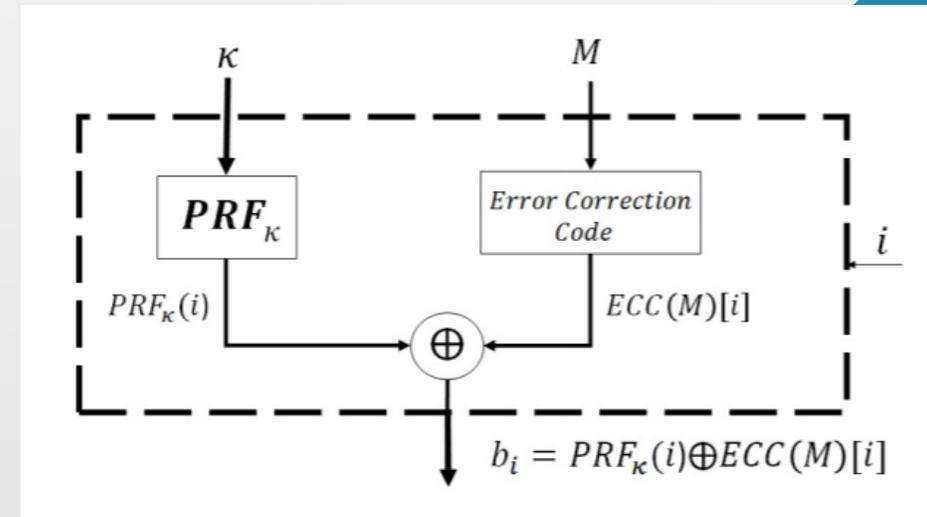
► **Problem [Indistinguishability]:** If the encoded message is $m = 00000000\dots$, the noise at the transitions will always be derived from P_{down}

► **Solution:** Sending pseudo-random bits, b_i , derived from m_i .

► $\kappa \xleftarrow{\$} \{0,1\}^l$ is a key, deployed at the sensor and actuator.

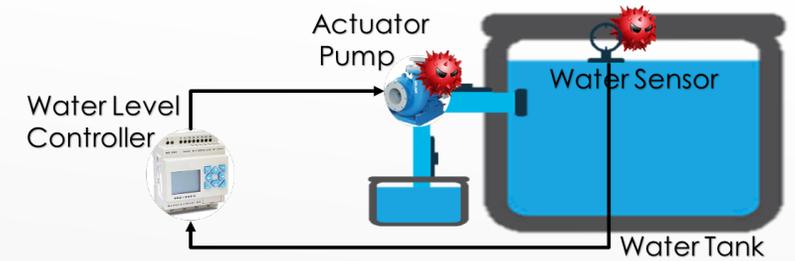
► **Assumption:** Sensor and actuator have a synchronized i .

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ATTACKER CHALLENGES

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Pseudo-random

Transition Parity

Evaluation

EVALUATION

- ▶ How good is the receiver in intercepting the chatty-sensor bits?
- ▶ **Theoretical:** Channel Capacity.
- ▶ **Practical:** Bit-error-rate of our chatty-sensor design.

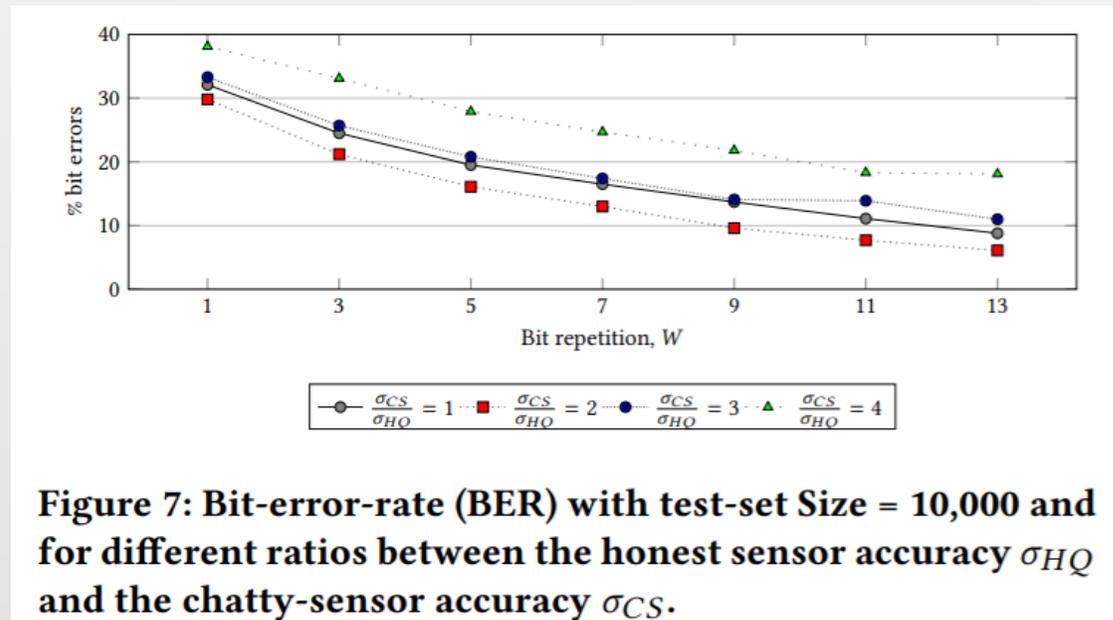
EVALUATION: CHANNEL CAPACITY

- ▶ **Channel Capacity** – highest information rate that can be achieved.
- ▶ Evaluated in a pH control process Simulink simulation.
- ▶ Based on real-world pH sensors noise.
 - ▶ σ_{CS}, σ_{HQ} - chatty-sensor / internal high-quality sensor noise standard deviation.
- ▶ **Results:** About 0.12 bit of information on every transition.
 - ▶ 1 transition every 5 seconds = 1.44 bits per minute.

$\frac{\sigma_{CS}}{\sigma_{HQ}}$	Avg. Bit Flip Probability	Avg. Channel Capacity
2	0.32	0.1
3	0.3	0.12
4	0.33	0.08
5	0.38	0.04

EVALUATION

- ▶ **Channel Capacity – 0.12 bit per transition.**
- ▶ **Bit-Error-Rate (BER)** – fraction of errors in the bits decoding.
 - ▶ Using repetition as error-correction-code: **~10% decoding errors**, with repetition of 13.
- ▶ We need better error-correction-codes for this channel [Future Work].



SUMMARY AND DISCUSSION

- ▶ Choosing devices based on **specification** and **price** enables **provable** covert attacks.
- ▶ As far as we know – this is the first **provable** covert channel from **sensors to actuator**.
- ▶ Requires to improve defenses:
 - ▶ Adding randomness to the channel (e.g. in the controller logic)
 - ▶ Purchasing devices from different vendors.
- ▶ In future works:
 - ▶ Improving the BER – maybe by non-provable method.

QUESTIONS?