

Securing Industrial Control Systems An E2E Integrity Verification Approach

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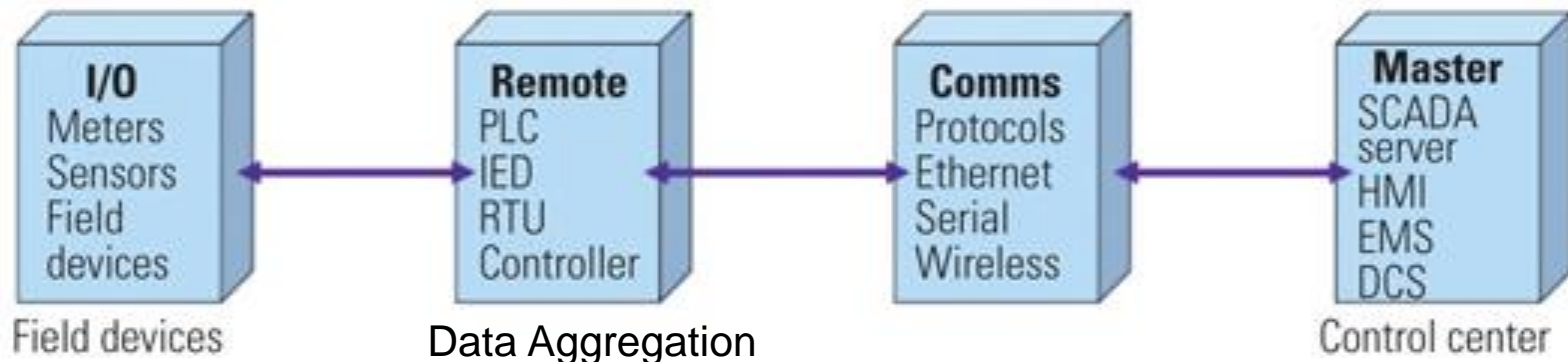
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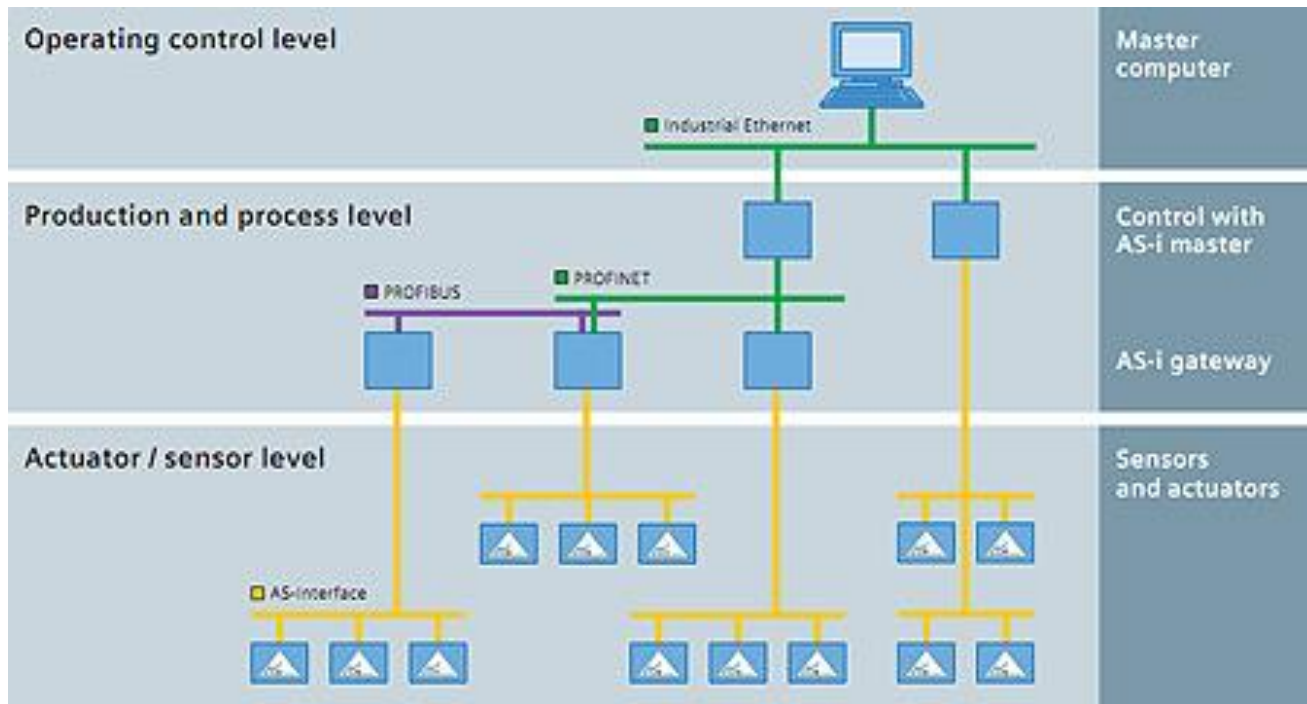


- Industrial Control Systems (ICS) are used to monitor and control industrial facilities and processes:
 - **Power Grid:** generation, distribution, load balancing and billing
 - **Chemical and Nuclear Plant:** control of safety critical processes.
 - **Gas and Water Facilities:** collect measurements from PLC/sensors and issue commands to actuators.

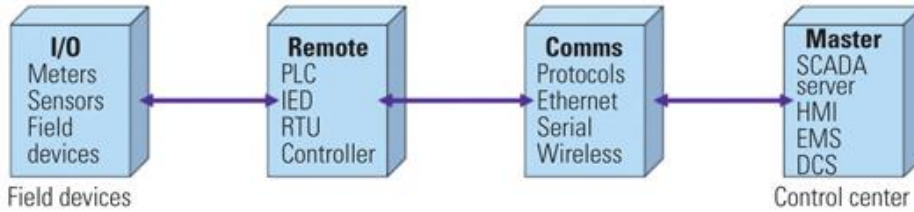


An Example ICS Architecture

- Master ensures data exchange with the slaves (field controller) by means of cyclic polling.
- Data collected at the field controller can be aggregated.



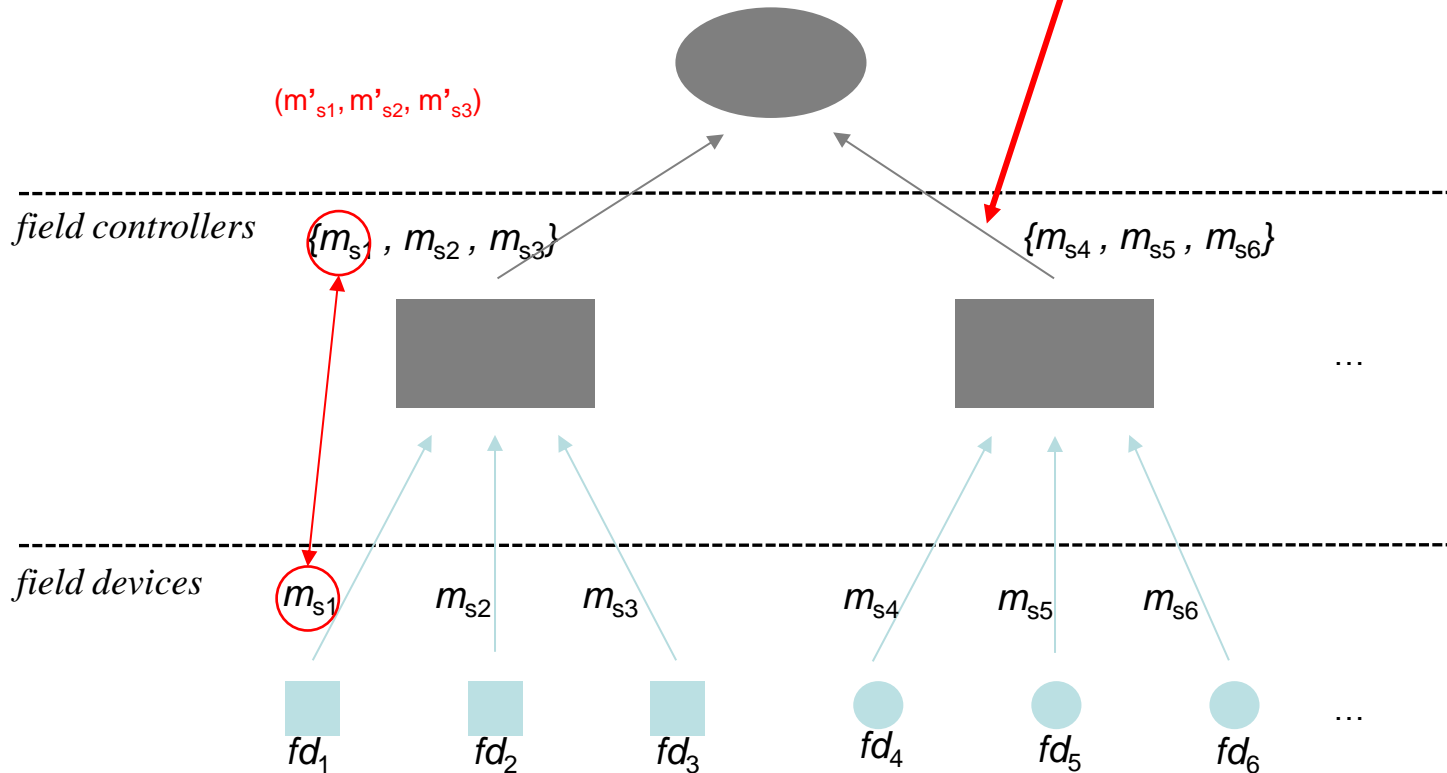
[Siemens]







Central controller

Vulnerabilities

- fraud
- selectively reporting
- single point of failure



- **Data Integrity** – the measurements on the field devices must reflect the current state of the instruments in the plant.
 modification and tampering.
- **Data Origin Authentication** – important to ensure that measurements are taken using the designated field devices.
 spoofing
- **Secure Data Aggregation** – though data are aggregated to save bandwidth, the central controller (Back End Master) must have the ability to check the integrity and data origin.
 integrity  data origin

- Chameleon Hashing
 - Hash function with a trapdoor for finding collusion.
 - Associated with a pair of public-private key.
 - Private-key serves as the trapdoor.
- Properties
 - Chameleon Hash Value [CHV] = $\text{CHA_HASH}(y, m, r)$.
 - given trapdoor x , find a collision $[m', r']$ where $m' \neq m$ and $r' \neq r$.
 - Hence $[\text{CHV}] = \text{CHA_HASH}(y, m', r')$.
- Chameleon Signature
 - Apply traditional signature, e.g., DSA, RSA, ECC to Chameleon Hash.



Field Devices



Trapdoor Hash Key
(x)



Trapdoor Chameleon Hash Function



Device ID
(Id_{fd})

Field Controllers



Chameleon Hash Key
(y)



Chameleon Hash Function



Chameleon Hash Key
(y)



Chameleon Hash Function

Back-end

Secure Channel

Secure Channel

Key Generation

- Krawczyk and Rabin's discrete logarithm construction
 - Two primes p and q are randomly generated such that $p = kq+1$ where q is a large prime factor.
- An element g of order q in \mathbb{Z}_p^* is chosen so that the private key, $x \in \mathbb{Z}_p^*$. The public-key, y is generated as

$$y = g^x \text{ mod } p$$

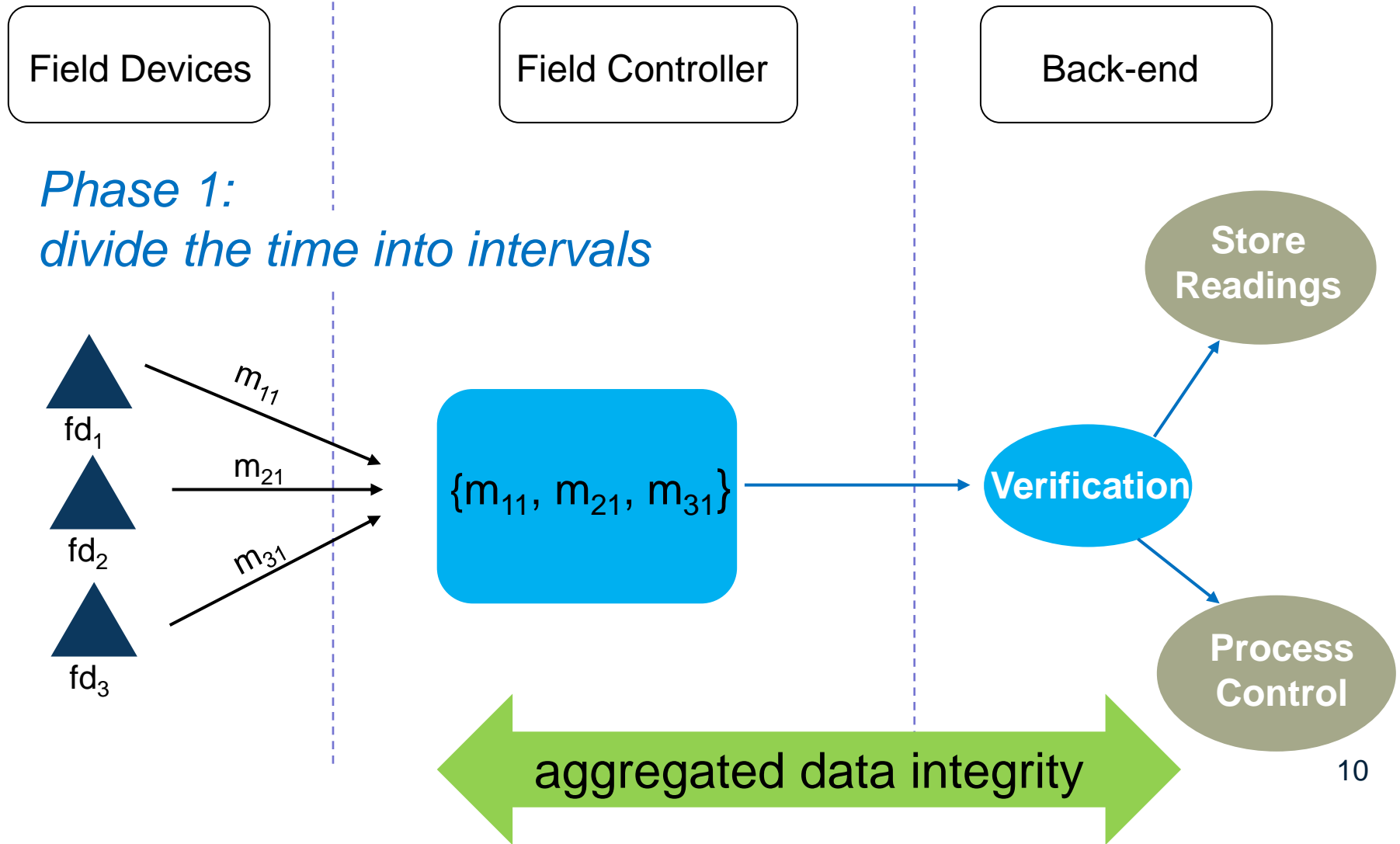
Generation of Chameleon Hash

- Given a message $m \in \mathbb{Z}_p^*$, choose a random value $r \in \mathbb{Z}_p^*$, the Chameleon Hash denoted as CHV can be computed as:

$$\text{CHA_Hash}(m,r) = g^m y^r \text{ mod } p$$

- Only the field devices have the ability to produce the same Chameleon Hash using a different message, m' such that $\text{CHA_Hash}(m,r) = \text{CHA_HASH}(m',r')$ by solving r'

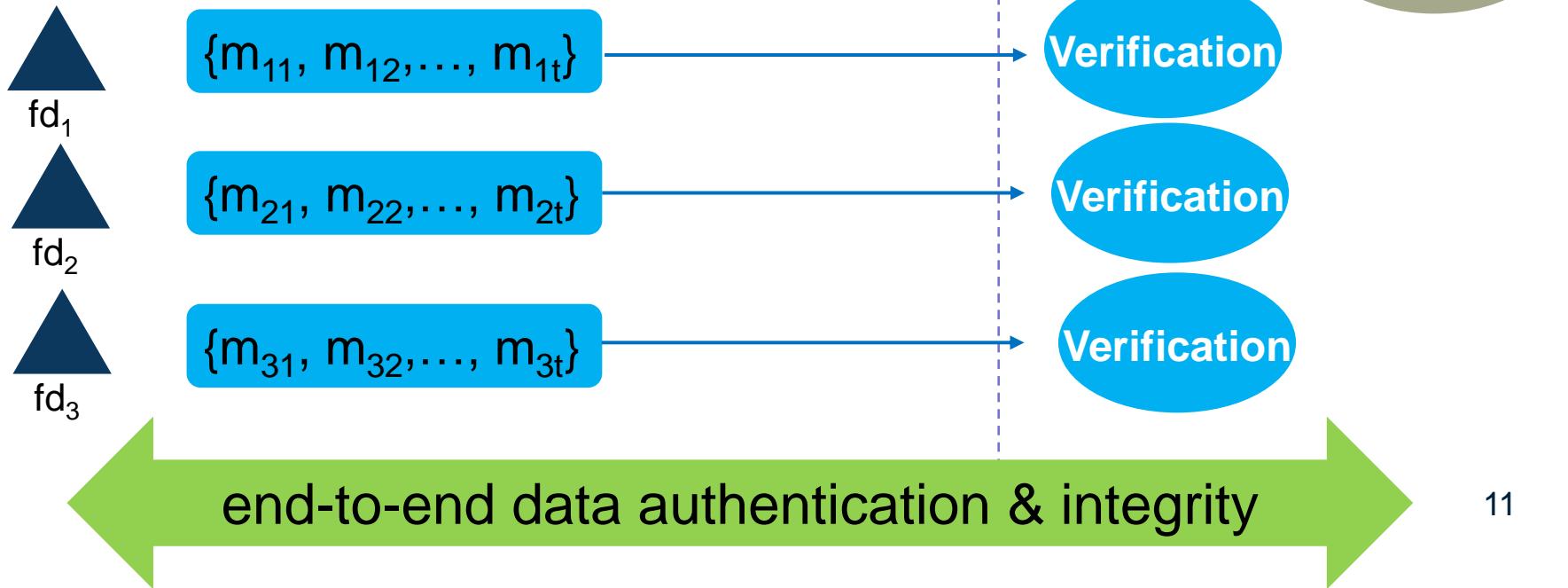
$$m + xr = m' + xr' \text{ mod } p$$

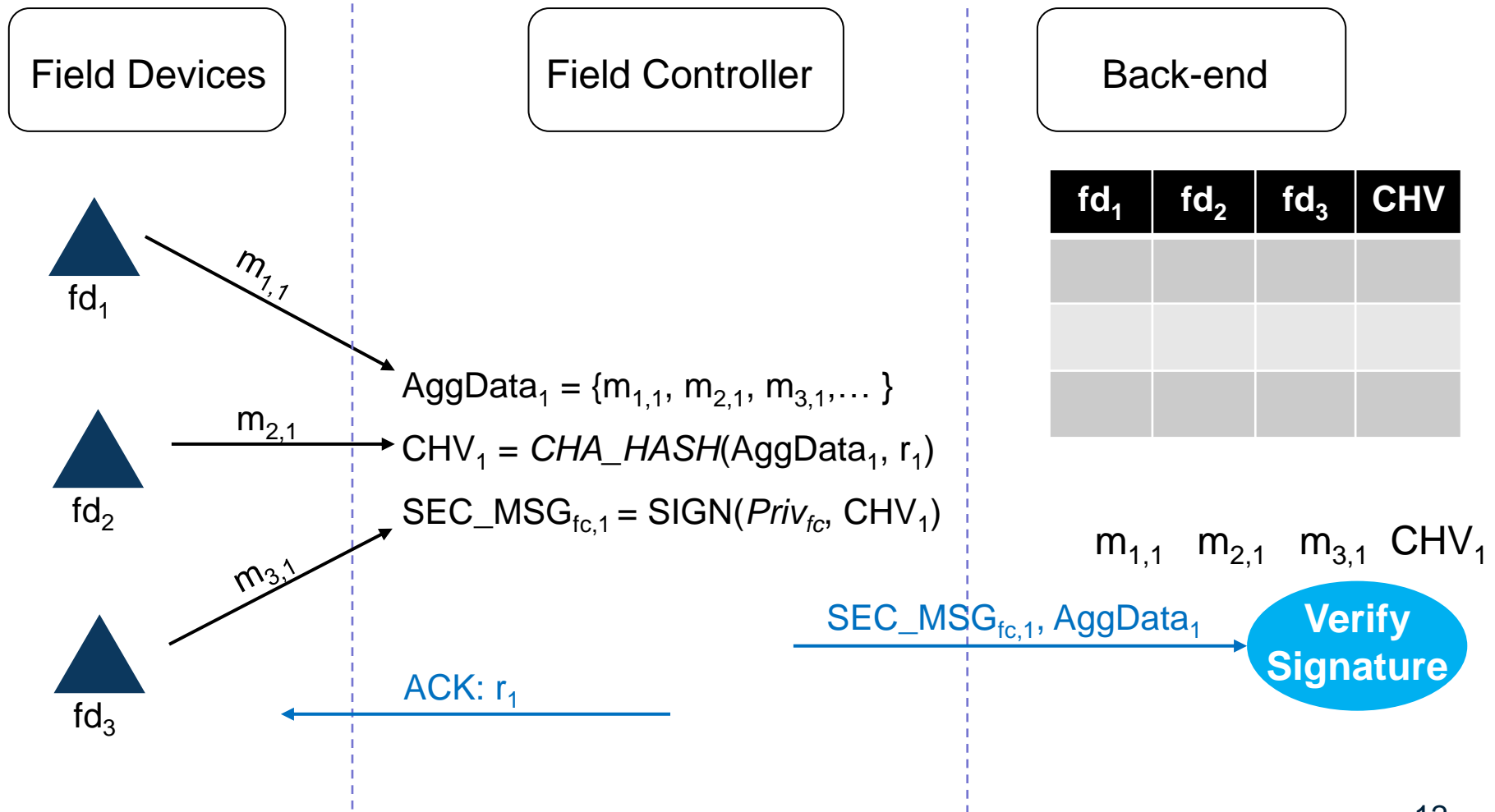


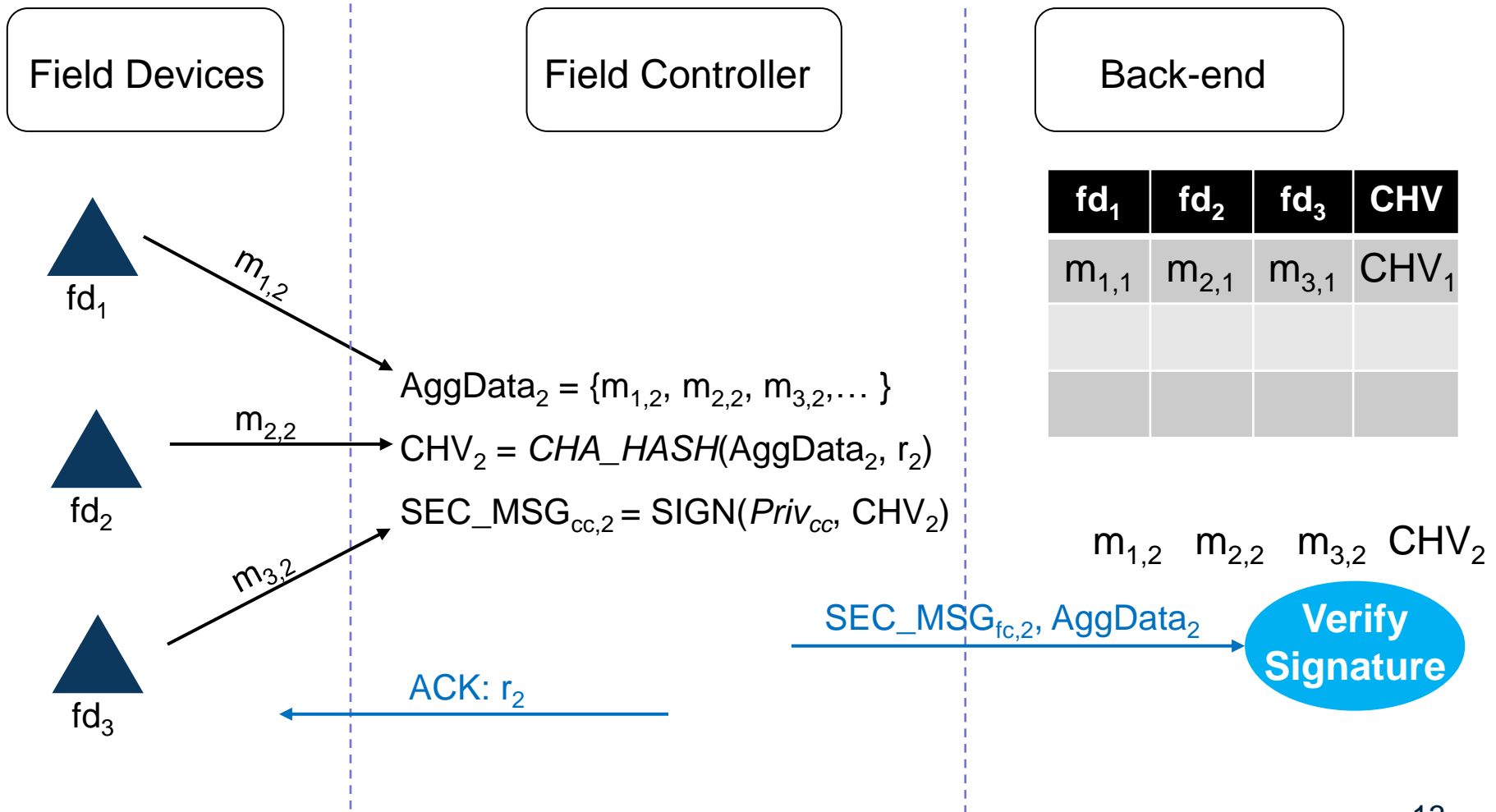
Field Devices

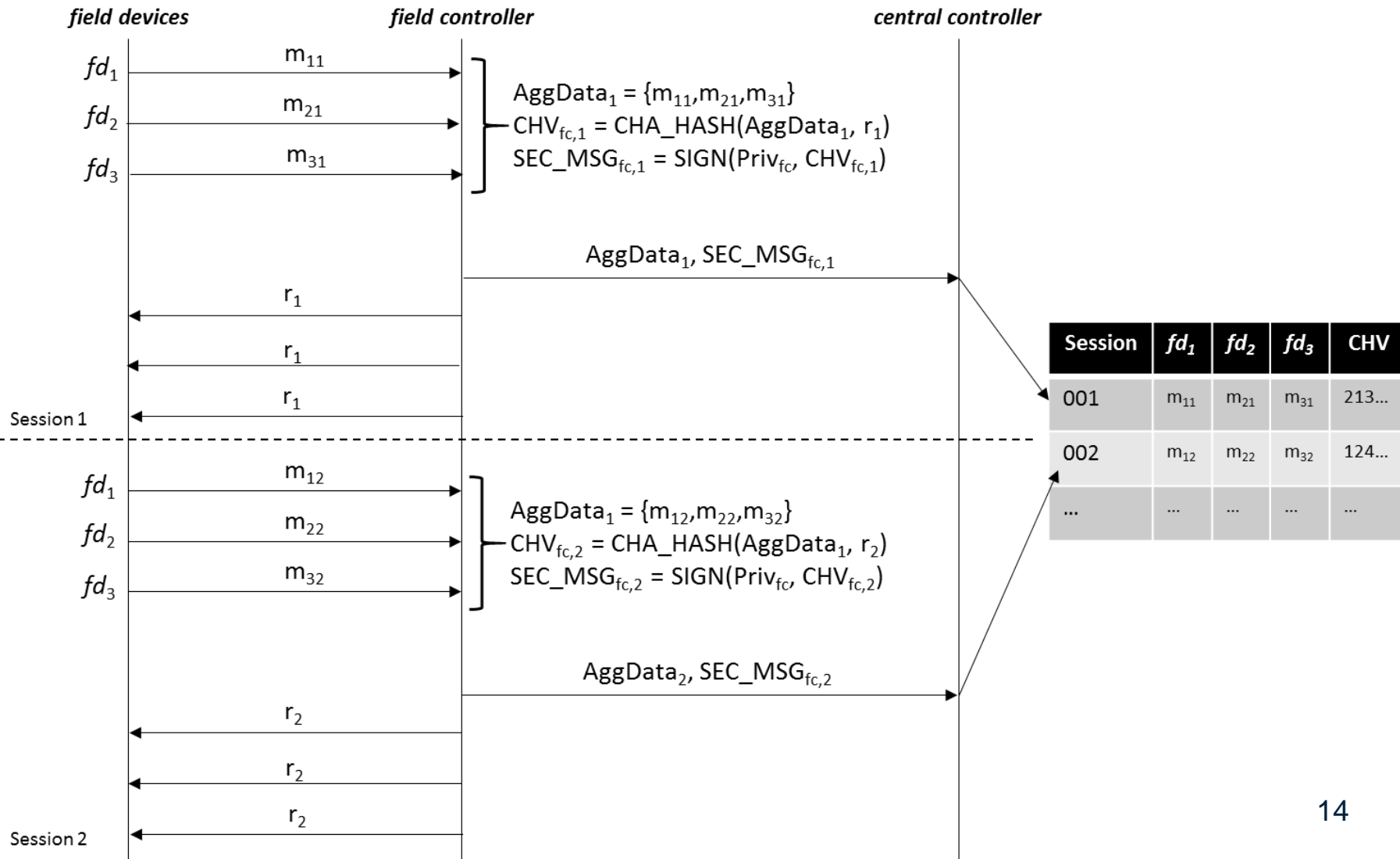
Back-end

Phase 2:
After t sessions in each interval









Transmission of Evidence

- Time is divided into intervals, where each interval consists of t sessions.
- At the end of each interval, field devices choose an r_v where $1 \leq v \leq t$, so that

$$\text{CHA_HASH}(m'_i, r'_i) = \text{CHA_HASH}(\text{AggData}_{v'}, r_{v'})$$

- m' denotes all the readings recorded by the field device i in the interval $\{Id_{fd,i}, m_{i,1}, m_{i,2}, \dots, m_{i,t}\}$

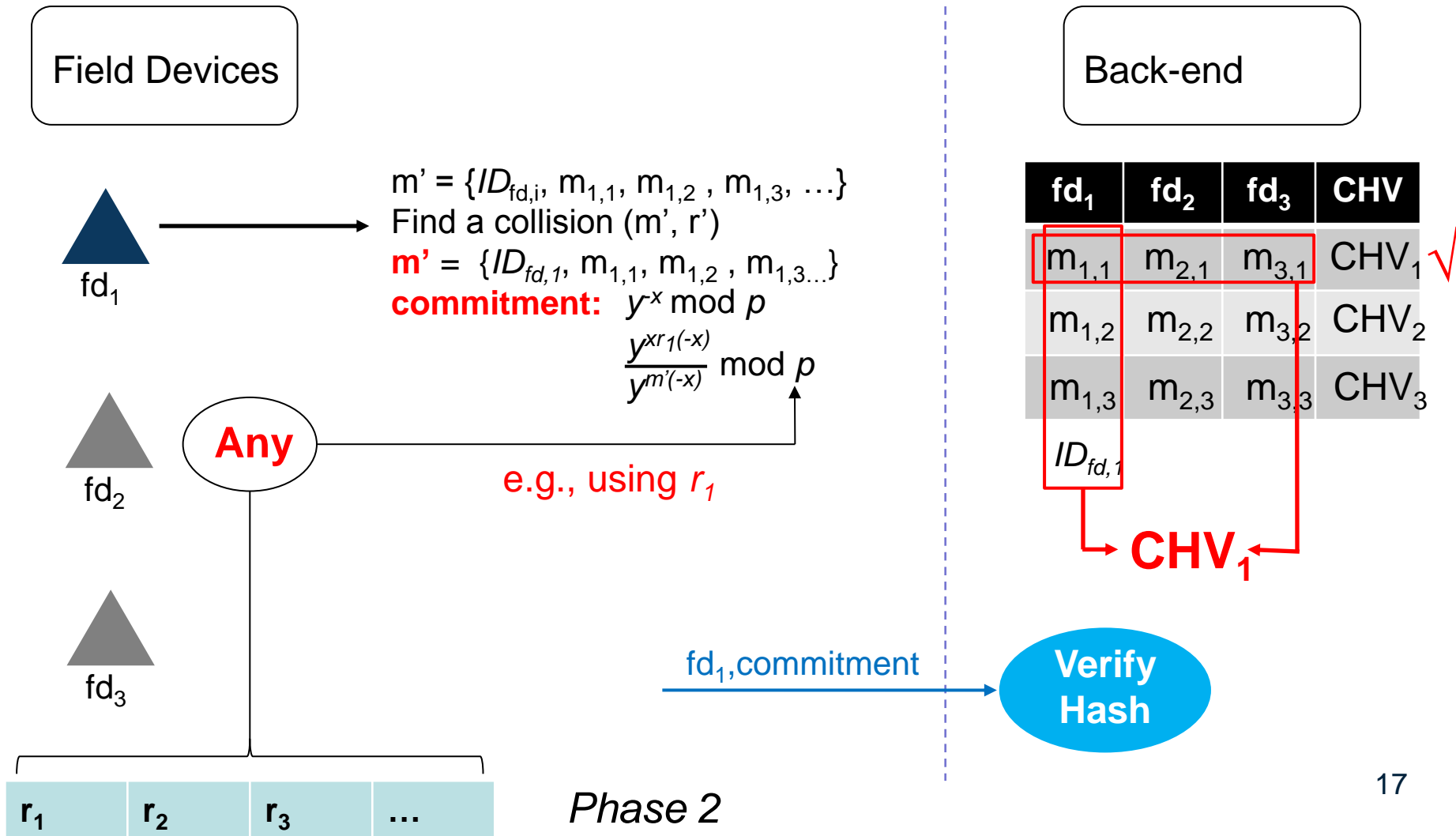
Transmission of Evidence

- To verify this, we need to solve r'_i

$$r'_i \bmod p = (\text{AggData}_v + xr_v - m')x^{-1} \bmod p$$

- However, field devices do not know AggData_v (sent by the field controller). Instead they can compute a **commitment** that allows the back-end to verify integrity and authenticity.

$$\left\{ y^x \bmod p \right\}, \left\{ \frac{y^{xr_v(-x)}}{y^{m'(-x)}} \bmod p \right\}$$

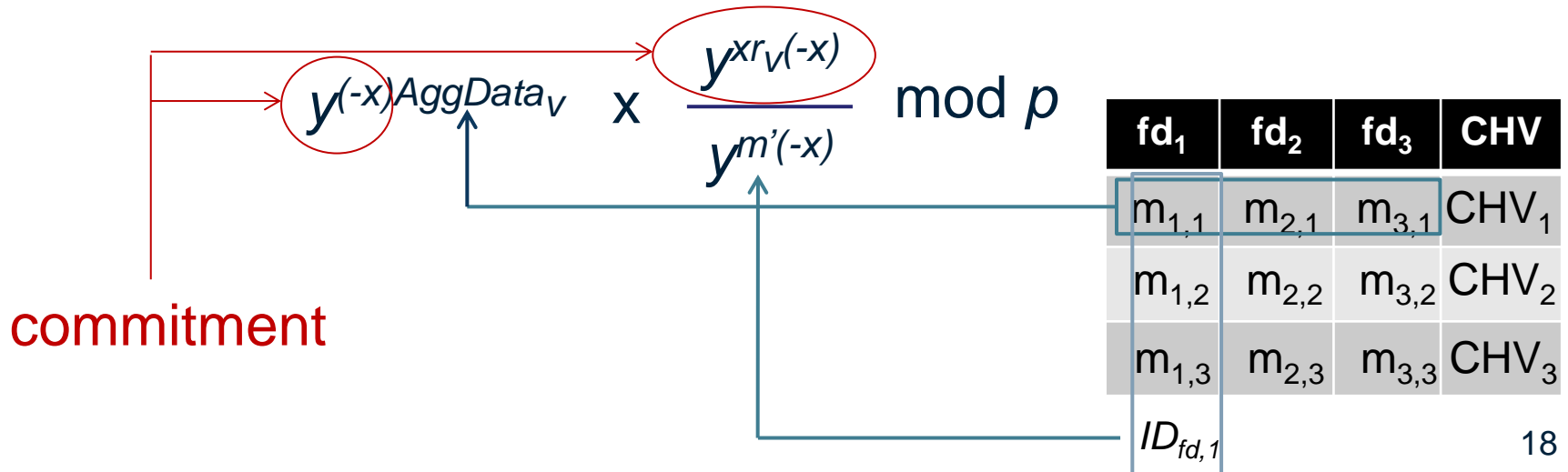


Integrity Verification

- We need to solve this:

$$r'_i \bmod p = (\text{AggData}_v + xr_v - m')x^{-1} \bmod p$$

- But, essentially we want to compute $\text{CHA_HASH}(m', r')$, so we need $y^{r'_i} \bmod p$, which is



- Prototype was implemented using Java, and deployed on Raspberry Pi Model B+
 - CPU: 700 MHz Low Power ARM processor
 - Memory: 512 MB
- Preliminary performance results

Device	Operation	Time (ms)
Controller	Chameleon Hashing	0.955955 (PC)
Field Device	Generation of Commitment	111.6 (Pi)
Back End	Integrity Verification	2.288591 (PC)
Field Device	Signature generation	5830 (Pi)

- Our scheme provides:
 - Data Integrity
 - Data Origin Authentication
 - Secure Data Aggregation
- Novel use of Chameleon Hashing and Signature other than its traditional usage, to detect **misbehaviour of controllers or aggregators in ICS/SCADA.**
- Future work:
 - Implement the protocol on real hardware or ICS platform.
 - Protocol can be generalized to be used in AMI, body sensor network, or any network with a hierarchical structure.

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